



Test-Adequacy

Test-suites Assessment Using Control Flow and Data Flow

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What is test adequacy?

It is necessary to know if the system has been tested thoroughly. The question is:

Is test suite T good enough?

Correspondingly this requires to define an **adequacy criterion** to make the assessment

Two different classes of criteria - to combine

- ▶ **Black-box**: based on models and requirements
- ▶ **White-box**: based on code

Example

Consider a program P developed to satisfy a set of requirements (P,R) (simplified version)

- **R1**: Input two integers, x, y , from the standard input device
- **R2**: Find and print to the standard output the sum if $x < y$
- **R3**: Find and print to the standard output the product of the two numbers if $x \geq y$
- **C**: A test T for program (P,R) is considered adequate if for each requirement r in R there is at least one test case in T that tests the correctness of P with respect to r

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Adequacy criteria push the improvements of test sets

Adequacy criteria are mainly meant as indicators to consider to **improve a test suite**.

- 1 Measure adequacy of T with respect to C . If T is adequate go to 3
- 2 For each uncovered element $e \in C_e$ do the following until e is covered or is determined to be infeasible
 - 1 construct a test t that covers e or will likely cover e
 - 2 execute P against t till no fault is identified
 - 3 if e is covered then t is added to T otherwise the tester can still decide to add it or to ignore it
- 3 The procedure ends

Example

A program computing x^y :

```
begin
  int x,y;
  int product, count;
  input(x,y);
  if (y >= 0) {
    product = 1; count = y;
    while (count > 0) {
      product = product * x;
      count = count - 1;
    }
    output(product);
  }
  else
    output("Input does not match its specification");
}
```

Criteria

- C1:** A test set is considered adequate if it tests the program for at least one zero and one nonzero value of each of the two inputs x and y
- C2:** A test set is considered adequate if it tests all paths. In case the program contains a loop, then it is adequate to traverse the loop body zero times and once.

Infeasibility

It is clearly possible that some criteria could be infeasible given P structure since some of the path are not possible

Checking infeasibility is not an easy task that in general cannot be decided. Is the tester that should decide if a path is feasible or not

```
begin
  int x,y;
  int z=0;
  input(x,y);
  if (x<0 and y<0) {
    z = x * x;
    if (y >= 0) {
      z = z + 1;
    } else {
      z = x * x * x;
    }
  }
  output(z);
}
end
```

Single or multiple executions

Warning

When a software **keeps state among different runs** it could be necessary to bring the system in a given state before being able to observe a failure

Criteria based on control flow

Statement coverage

The statement coverage of T with respect to (P,R) is computed as $|S_c|/(|S_e| - |S_i|)$ where S_c is the set of statements covered, S_i the set of unreachable statements, and S_e the set of statements in the program, that is the coverage domain. T is considered adequate with respect to the statement coverage criterion if the **statement coverage of T with respect to (P,R) is 1**.

Block coverage

The block coverage of T with respect to (P,R) is computed as $|B_c|/(|B_e| - |B_i|)$ where B_c is the set of blocks covered, B_i the set of unreachable blocks, and B_e the blocks in the program, that is the block coverage domain. T is considered adequate with respect to the block coverage criterion if **the block coverage of T with respect to (P,R) is 1**.

Conditions and decisions

- Conditions can be classified as **simple** or **compound**
- Conditions are generally used to define **decision points**
- A **decision is covered** if the flow has been diverted to all possible destinations

Decision Coverage

The decision coverage of T with respect to (P,R) is computed as $|D_c|/(|D_e| - |D_i|)$ where D_c is the set of decisions covered, D_i the set of unfeasible decision, and D_e the set of decision in the program, that is the decision coverage domain. T is considered adequate with respect to the decision coverage criterion if **the decision coverage of T with respect to (P,R) is 1**.

To be considered are peculiarities related to the `switch` statements

Condition Coverage

The condition coverage of T with respect to (P,R) is computed as $|C_c|/(|C_e| - |C_i|)$ where C_c is the set of simple conditions covered, D_i the set of unfeasible simple conditions, and C_e is the set of simple conditions in the program, that is the condition coverage domain. T is considered adequate with respect to the decision coverage criterion if **the decision coverage of T with respect to (P,R) is 1**.

Condition vs. decision coverage

Condition coverage does not guarantee decision coverage and viceversa

Condition/decision coverage

The condition/decision coverage of T with respect to (P,R) is computed as $(|C_c| + |D_c|) / ((|C_e| - |C_i|) + (|D_e| - |D_i|))$ where variable as defined as before. T is considered adequate with respect to the condition/decision coverage criterion if **the condition/decision coverage of T with respect to (P,R) is 1.**

Example

Consider a program that takes in input two integers x and y , and returns an integer z according to the following table:

$x < 0$	$y < 0$	output(z)
true	true	foo1(x, y)
true	false	foo2(x, y)
false	true	foo2(x, y)
false	false	foo1(x, y)

Apply the test suite $T = \{t_1 : \langle x = -3, y = -2 \rangle, t_2 : \langle x = -4, y = 2 \rangle\}$ to the program below

```
begin
  int x, y, z;
  input(x, y);
  if (x < 0 and y < 0)
    z = foo1(x, y);
  else
    z = foo2(x, y);
  output(z);
end
```

Multiple Condition Coverage

This criterion aims at assessing the software with **all possible combinations of simple conditions** constituting a compound condition

Multiple condition coverage

The multiple condition coverage of T with respect to (P,R) is computed as $|C_c|/(|C_e| - |C_i|)$ where $|C_c|$ denotes the set of combinations covered, $|C_i|$ denotes the set of infeasible simple combinations, and $|C_e|$ is the total number of combinations in the program. T is considered adequate with respect to the multiple-condition coverage criterion if **the multiple-condition coverage of T with respect to (P,R) is 1.**

Let's consider a code composed of n decisions each one including K_i with $i \in [1 \dots n]$ simple conditions. In case all of them are feasible which is the total number of possible combinations?

Example

Consider a program that takes in input three integers A , B and C , and returns a value S according to the following table:

$A < B$	$A > C$	S
true	true	$f_1(A, B, C)$
true	false	$f_2(A, B, C)$
false	true	$f_3(A, B, C)$
false	false	$f_4(A, B, C)$

Apply the test suite $T = \{t_1 : \langle A = 2, B = 3, C = 1 \rangle, t_2 : \langle A = 2, B = 1, C = 3 \rangle\}$ to the program below

```
1 begin
2   int A, B, C, S=0;
3   input (A, B, C);
4   if (A<B and A>C) S=f1(A, B, C);
5   if (A<B and A<=C) S=f2(A, B, C);
6   if (A>=B and A<=C) S=f4(A, B, C);
7   output (S);
8 end
```

Modified Condition/Decision Coverage – MC/DC

- ▶ Combinations necessary to satisfy the **Multiple Condition Coverage** is generally too big.
- ▶ MC/DC allows a coverage of all decisions and all conditions avoiding the exponential explosion
- ▶ To derive the test set the idea is to identify those tuples which can cover the two criteria without requiring a complete combinations of values.

Let's consider the compound condition $(C_1 \wedge C_2) \vee C_3$

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Definition of MC/DC coverage

The MC/DC criterion requires that:

- Each **block** in P has been covered
- Each **simple condition** in P has taken both `true` and `false` value
- Each **decision** in P has taken all possible outcomes
- Each simple condition within a compound condition C in P has been shown to independently affect the outcome of C (**limited to the simple condition when it occurs more than once**).

Measure

Measure the 4 different factors separately and for MC:

$$\blacktriangleright MC_C = \frac{\sum_{i=1}^N e_i}{\sum_{i=1}^N (n_i - f_i)}$$

where n_i number of simple conditions, e_i single conditions for which independent effects have been shown, f_i number of infeasible conditions.

MC/DC comparisons

MC/DC vs. Condition

- ▶ Missing conditions
- ▶ Incorrect boolean operator
- ▶ Mixed type

MC/DC vs. Multiple time comparison

<i>n</i>	Multiple Condition	MC/DC	Multiple Condition	MC/DC
1	2	2	2ms	2ms
4	16	5	16ms	5ms
8	256	9	256ms	9ms
16	65536	17	65.6s	17ms
32	4294967296	33	49.5 days	33ms

MC/DC and Lazy evaluation

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MC/DC and Lazy evaluation

Example

Consider a program conceived to satisfy the following requirements:

R_1 : Given coordinate position x , y , and z , and a direction value d , the program must invoke one of the three functions `fire-1`, `fire-2`, and `fire-3` as per conditions below:

$R_{1,1}$: Invoke `fire-1` when $(x < y)$ and $(z * z > y)$ and (`prev`="East") where *prev* and *current* denote, respectively, the previous and current values of d .

$R_{1,2}$: Invoke `fire-2` when $(x < y)$ and $(z * z \leq y)$ or (`current`="South")

$R_{1,3}$: Invoke `fire-3` when none of the two conditions above is `true`

R_2 : The invocation described above must continue until an input Boolean variable becomes `true`

- ▶ let's generate test satisfying the conditions and let's analyze the covered decision on a possible implementation of the system

Code

```
begin
float x,y,z; direction d; string prev,current; bool done;
input(done); current ='North';
while(!done) {
    input(d); prev=current;current=f(d); input(x,y,z);
    if ((x<y) and (z*z>y) and (prev=='East'))
        fire-1(x,y);
    else if ((x<y) and (z*z <= y) or (current == 'South'))
        fire-2(x,y);
    else
        fire-3(x,y); input(done);
}
output('Firing completed');
end
```

- ▶ generate tests to meet the requirements (4 tests generated)

Test	Req.	done	d	x	y	z
t_1	$R_{1,2}$	false	East	10	15	3
t_2	$R_{1,1}$	false	South	10	15	4
t_3	$R_{1,3}$	false	North	10	15	5
t_4	R_2	true				

- ▶ Which kind of coverage criteria are satisfied by the test set?
 - ▶ Cover $x < y$ to get condition coverage?
 - ▶ What about Multiple Condition Coverage?
 - ▶ What about MC/DC?

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Tracing test cases to requirements

Enhancing a test set we should understand *what portions of the requirements are tested when the program under test is executed against the newly added test case?*

- Trace back test to requirements is useful when they need to be modified

Data Flow concepts

- Criteria considered so far are based on the **control flow**
- it is possible to conceive adequacy criteria based on **data flow characteristics**

Consider the following program:

```
begin
  int x,y; float z;
  input(x,y);
  z=0;
  if (x!=0) z=z+y;
    else z=z-y;
  if (y!=0) z=z/x // Should be (y!=0 and x!=0)
    else z=z*x;
  output(z);
end
```

An MC/DC test set could not reveal the error while a test set based on definition and usage of variables would have been able

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Data flow criteria

- ▶ Data flow criteria based on two main concepts:
 - **Definitions** – points in which a variable is defined (e.g. assignments, input statements)
 - **Uses** – points in which a variable is accessed
 - **computational usage** - c-use
 - **predicate usage** - p-use
- ▶ Which are the effect of parameter passing (**by value or by reference**)?

```
input (x, y); z=0;
z = x+1
A[x-1]=B[2];
foo(x*x);
output(z);
if (z>0) output(x);
if (A[x+1]>0) output(x);
```

Global, Local and Pointers

Variables can be defined in a block, used and redefined (**killed**) within the same block. Effects can also be available outside the block:

```
▶ p = y+z; x = p+1; p = z*z;
```

Definition and use of variables can be referred to:

- local
- global

Pointers

The **use of pointers** makes the data flow analysis rather complex:

```
z=&x;  
y=z+1;  
*z=25;  
y=*z+1;
```


Data Flow Graph

A **data-flow** graph of a program (aka def-use graph) captures the flow of definitions across the basic blocks constituting the program. The graph can be constructed in the following way:

- 1 Construct def_i , $c - use_i$, $p - use_i$ for each basic block i in P
- 2 Associate each node i in N with def_i , $c - use_i$, $p - use_i$
- 3 For each node i that has a non empty $p - use$ set and ends in condition C , associate edges (i, j) and (i, k) with C and $!C$, respectively.

Data flow graph

Build the DFG for the following piece of code:

```
begin
  int x,y,z;
  input(x,y); z=0;
  if (x<0 and y<0) {
    z=x*x;
    if (y>=0) z=z+1;
  }
  else z=x*x*x;
  output(z);
end
```

Example

Let's build a def-use graph for the following program:

```
begin
  float x,y,z=0.0; int count; input (x,y,count);
  do {
    if (x<=0) {
      if (y>= 0 {
        z=y*z+1;
      }
    } else { z= 1/x; }
    y=x*y+z; count = count -1;
  } while (count > 0)
  output (z);
end
```

Which are the blocks?

Example

Let's build a def-use graph for the following program:

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    } else { z= 1/x; }
    y=x*y+z; count = count -1;
  } while (count > 0)
  output (z);
end
```

Which are the blocks?

Definitions

def-clear paths

A def-clear path for a variable x is a path from a definition of the variable to a usage **without further definitions** in the intermediate nodes of the path

live definition

A definition at node i **is live** at node j if there is no intermediate definition in a path from i to j

Def-use pairs

A **def-use pair** for a variable 'X' refers to a definition d and a usage u on a def-clear path

For each variable definition $d_i(x)$ there is:

- a **dcu**($d_i(x)$) set, that is constituted by all nodes j in any def-clear path from node i such that $u_j(x)$ in relation to a c-use
- a **dpu**($d_i(x)$) set, that is constituted by all sets of edges leaving a node j for which there is a def-clear path from i and $u_j(x)$ in relation to a p-use

Let's fill the table for the previous DFG

Variable(v)	Defined at node(n)	dcu(v,n)	dpu(v,n)

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Let's fill the table for the previous DFG

Variable(v)	Defined at node(n)	dcu(v,n)	dpu(v,n)

Def-use chains

A **def-use chain** (aka **k-dr interaction**) is constituted by path including a sequence of alternating def-use pairs. It is also possible to consider **different variables** E.g. consider the def-use chain for y and z and the sequence of nodes in which they are considered

Def-use optimization

Some def-use can subsume other relations. E.g. z defined in node 1 subsumes y defined in node 1

Let's fill the table for the previous DFG

Variable(v)	Defined at node(n)	$dcu(v,n)$	$dpu(v,n)$

Adequacy criteria for data-flow

Given the total number of c-uses (CU) and p-uses (PU) for all variable definitions we can define different coverage criteria for data-flow.

$$CU = \sum_{i=1}^n \sum_{j=1}^{d_i} |\mathbf{dcu}(v_i, n_j)|$$

$$PU = \sum_{i=1}^n \sum_{j=1}^{d_i} |\mathbf{dpu}(v_i, n_j)|$$

where $v = \{v_1, v_2, \dots, v_n\}$ is the set of variables in a program and $n = \{n_1, n_2, \dots, n_k\}$ is the set of blocks in the same program

Coverage

C-use coverage

The c-use coverage of T with respect to (P,R) is computed as:

$$\frac{CU_c}{CU - CU_f}$$

where CU_c is the number of c-uses covered and CU_f the number of infeasible c-uses. T is considered adequate with respect to the c-use coverage criterion if its c-use coverage is 1.

P-use coverage

The p-use coverage of T with respect to (P,R) is computed as:

$$\frac{PU_c}{PU - PU_f}$$

where PU_c is the number of p-uses covered and PU_f the number of infeasible p-uses. T is considered adequate with respect to the p-use coverage criterion if its p-use coverage is 1.

Coverage

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Coverage's

All-uses coverage

The all-uses coverage of T with respect to (P,R) is computed as:

$$\frac{CU_c + PU_c}{(CU + PU) - (CU_f + PU_f)}$$

where CU_c and PU_c are the number of c-uses and p-uses covered respectively. CU_f and PU_f are the number of infeasible c-uses and p-uses respectively. T is considered adequate with respect to the all-uses coverage criterion if its all-uses coverage is 1.

k-dr chain coverage

For a given $K \geq 2$ the $kdr(k)$ coverage of T with respect to (P,R) is computed as:

$$\frac{C_c^k}{C^k - C_f^k}$$

where C_c^k is the number of k-dr interactions covered, C^k is the number of elements in $K\text{-dr}(k)$, and C_f^k the number of infeasible interactions in $k\text{-dr}(k)$. T is considered adequate with respect to the $kdr(k)$ coverage criterion if its $k\text{-dr}(k)$ coverage is 1.

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All-uses coverage

The all-uses coverage of T with respect to (P,R) is computed as:

$$\frac{CU_c + PU_c}{(CU + PU) - (CU_f + PU_f)}$$

where CU_c and PU_c are the number of c-uses and p-uses covered respectively. CU_f and PU_f are the number of infeasible c-uses and p-uses respectively. T is considered adequate with respect to the all-uses coverage criterion if its all-uses coverage is 1.

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$$\frac{C_c^k}{C^k - C_f^k}$$

where C_c^k is the number of k-dr interactions covered, C^k is the number of elements in $K-dr(k)$, and C_f^k the number of infeasible interactions in $k.dr(k)$. T is considered adequate with respect to the $kdr(k)$ coverage criterion if its $k-dr(k)$ coverage is 1.

Control flow vs. Data Flow

The subsumes relation

A coverage criterion C1 subsumes a coverage criterion C2 iff whenever the satisfaction of C1 implies the satisfaction of C2

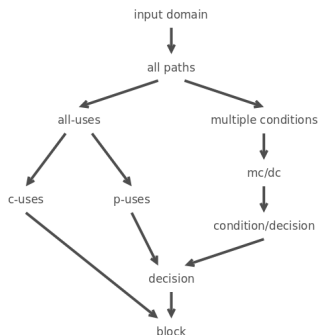


Figure: The subsumes relationship among the studied coverage criterion

Mutation analysis - Ch. 8

Sketch of the idea

Mutation is a powerful strategy to assess the **quality of test suites**. The approach is based on the generation of **program mutants** and on the score got by a test suite in “**killing**” them.

Regression testing - Ch. 9

Sketch of the idea

Definition of strategies to select subset of test cases in a test suite in order to test a system that has undergone a modification in order to reduce the costs of testing obviously getting enough confidence on the quality of the software.