
# Petri Nets

Business Process Management and Flexibility

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# Carl Adam Petri

- Introduced in 1962 (Petri's PhD thesis)
- 60's and 70's main focus on theory 80's focus on tools and applications
- Now applied in several fields
- Success due to simple and clean graphical and conceptual representation

(1926 – 2010)

# Formal and abstract business process specification

- **Formal**: the semantics of process becomes well defined and not ambiguous
- **Abstract**: execution environment is disregarded (Remind about separation of concerns)

- Petri Net can be used to model dynamic system with a static structure:
  - The static structure is represented by a Petri Net (process level)
  - The dynamic behavior is capture by the token play of the Petri Net (instance level)

# Basic Elements

## Place

- A place can stand for a state a medium a buffer a condition a repository of resources a type ...

## Transition

- A transition can stand for an event an operation a transformation a transportation a task an activity ...

## Token

- A token can stand for a physical object a piece of data a resource an activation mark a message a document a case ...

# Petri Net - Definition

A Petri Net is a tuple $(P, T, F, M_0)$ where

- P is a finite set of places
- T is a finite set of transitions
- $F \subseteq (P \times T) \cup (T \times U)$ is a flow relation
- $M : P \rightarrow N$ is the initial marking

# Marking

- A Marking M: P → N denotes the number of tokens in each place
- The marking of a Petri net represents its state
- *M (a) = 0* denotes the absence of tokens in a place *a*

# Firing a trasition

- Let (P, T, F) be a Petri net and M a marking.
- The firing of a transition is represented by a state change of the Petri net
  - $M \rightarrow M'$ indicates that by firing t, the state of the Petri net changes from M to M'
  - $M \rightarrow M'$ indicates that there is a transition t such that $M \rightarrow M'$
  - $M_1 \overset{*}{\rightarrow} M_n$ means that there is a sequence of transition $t_1, t_2, \ldots, t_{n-1}$ *for* $1 <= i < n$
  - A state M' is reachable from a state M if and only if $M \overset{*}{\rightarrow} M'$

# Pre-set and Post-set - Transitions

A place $p$ is an input place for trantition $t$ iff
$(p, t) \in F$
We let $\bullet t$ denote the set of input places of $t$.
(pre-set of $t$)

A place $p$ is an output place for trantition $t$ iff
$(t, p) \in F$
We let $t\bullet$ denote the set of output places of $t$.
(post-set of $t$)

# Pre-set and Post-set - Place

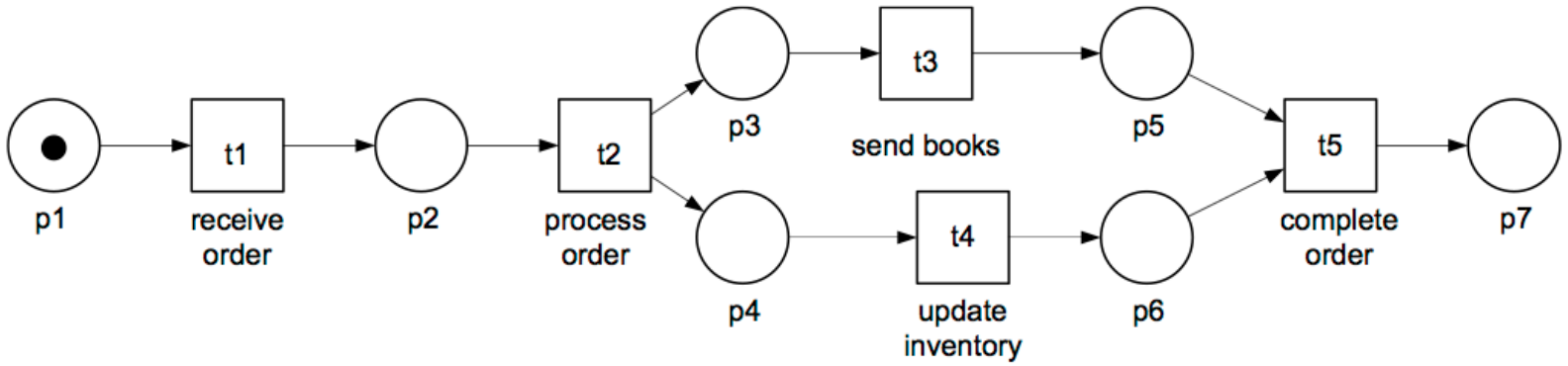
Analogously, we let

- $\bullet p$ denote the set of transitions that share $p$ as output place
- $p\bullet$ denote the set of transitions that share $p$ as input place

Formally:

$$\bullet x = \{\ y \mid (y, x) \in F\ \}$$
$$x\bullet = \{\ y \mid (x, y) \in F\ \}$$

# Sample petri net – single instance






GIVE THE FORMAL DEFINITION OF THE PETRI
in term of set T, P, F, pre-set e post-set!!!

# A Petri Net is a tuple $(P, T, F, M_0)$ where

- P is a finite set of places = (p1, p2, p3 …. p7)
- T is a finite set of transitions = (t1, …. t5)
- $F \subseteq (P \times T) \cup (T \times U)$ is a flow relation F= ((p1,t1) (t1, p2) (p2, t2) … )
- $M : P \rightarrow N$ is the initial marking M = (1, 0, 0 …. 0)

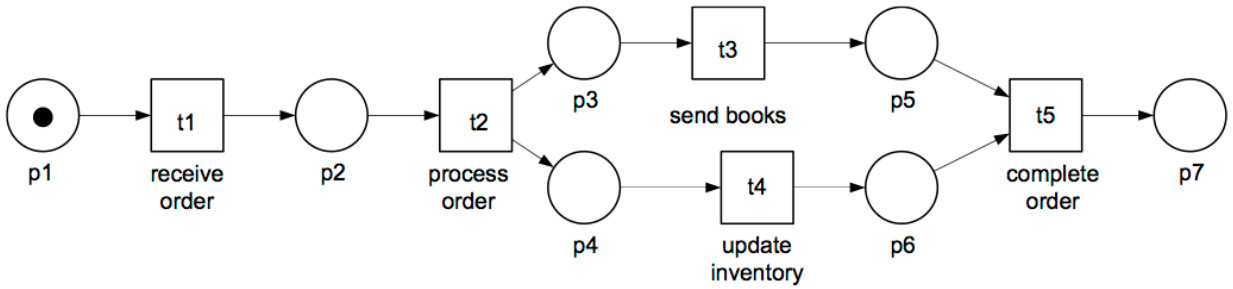
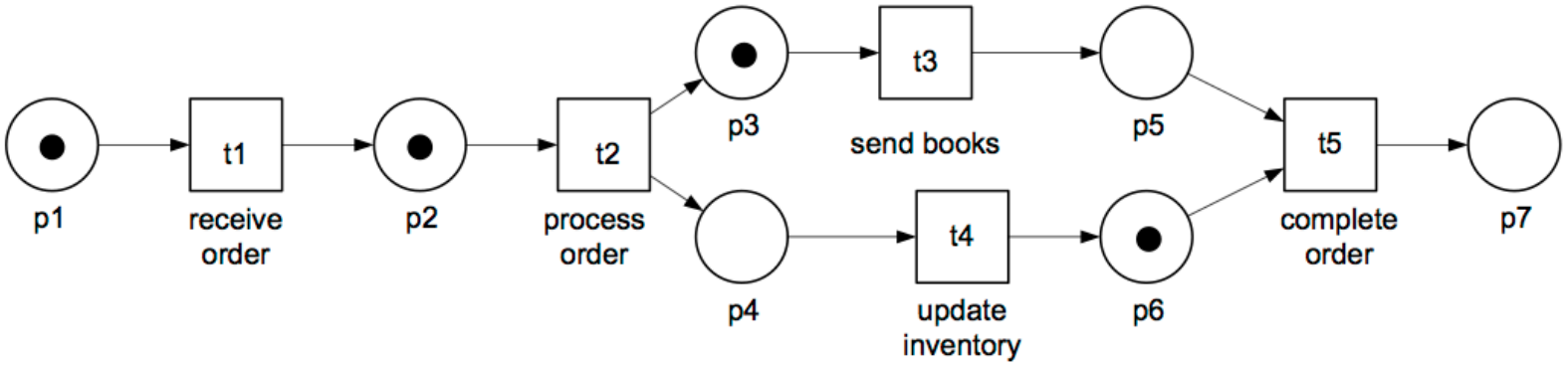



**Fig. 4.26.** Sample Petri net representing single process instance

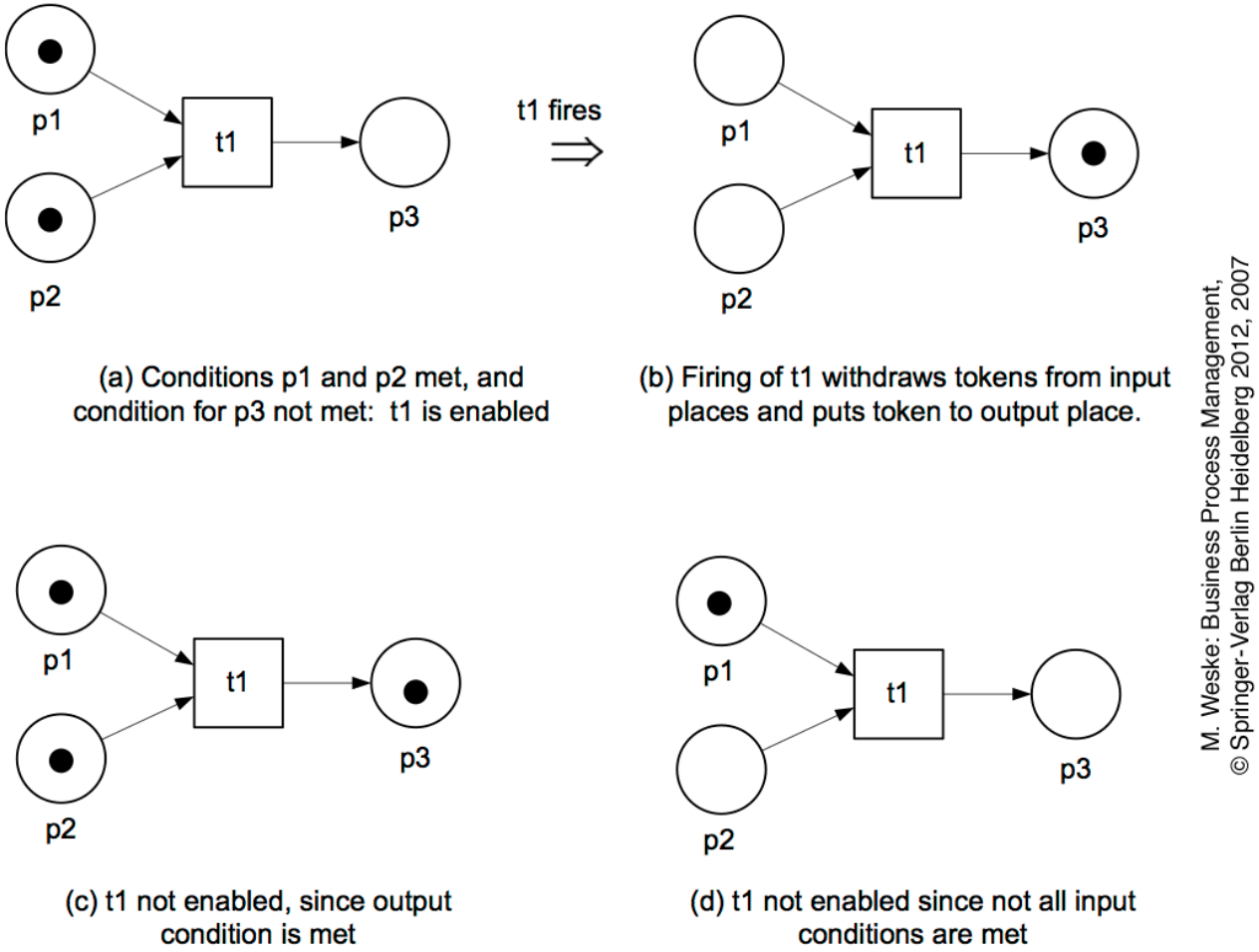
# Sample petri net – multiple instance



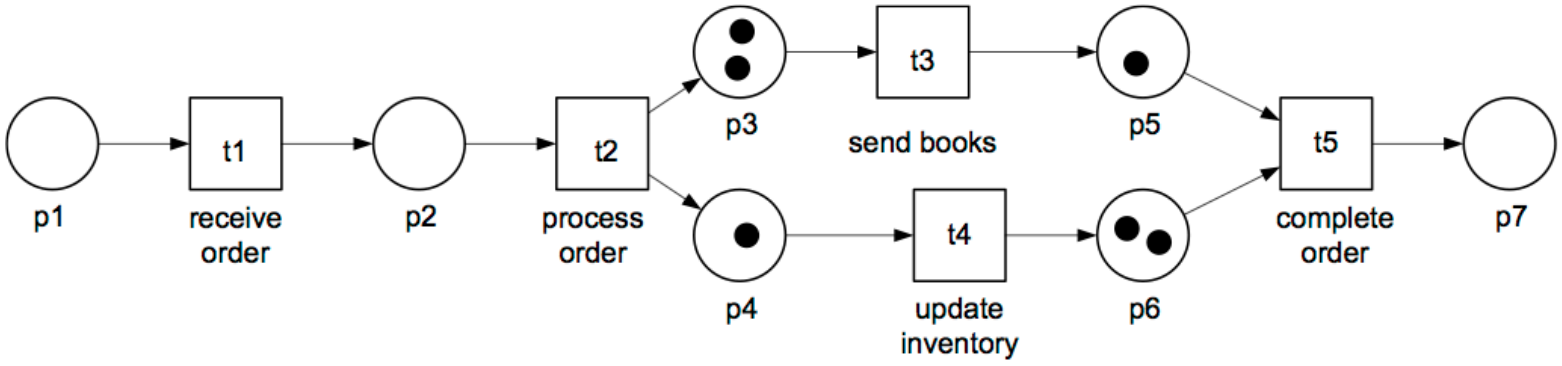



What are the instance active?
Give the marking for earch active instance!

# Safe-petri net or 1-bounded



# Sample Petri net – multiple instance

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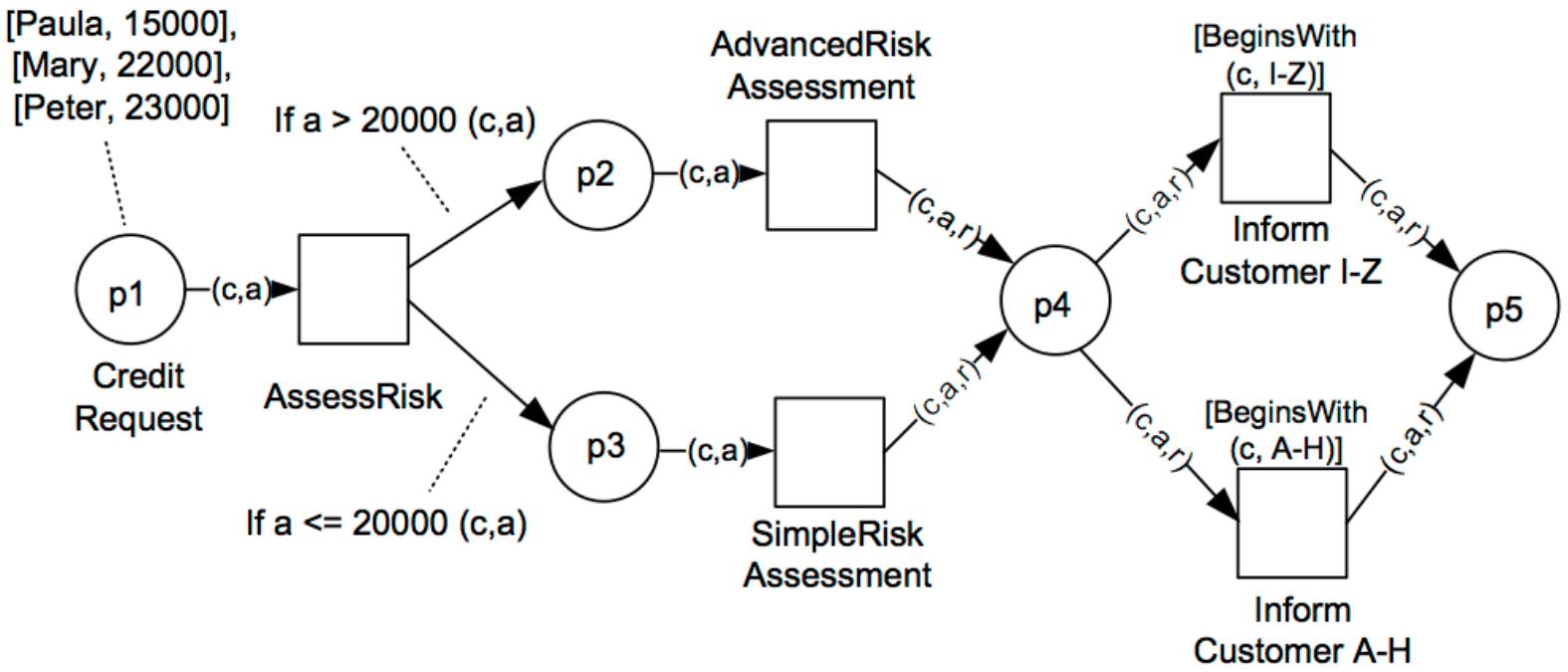
GIVE THE MARKING FOR EACH INSTANCE!!!!

# Colured Petri Net

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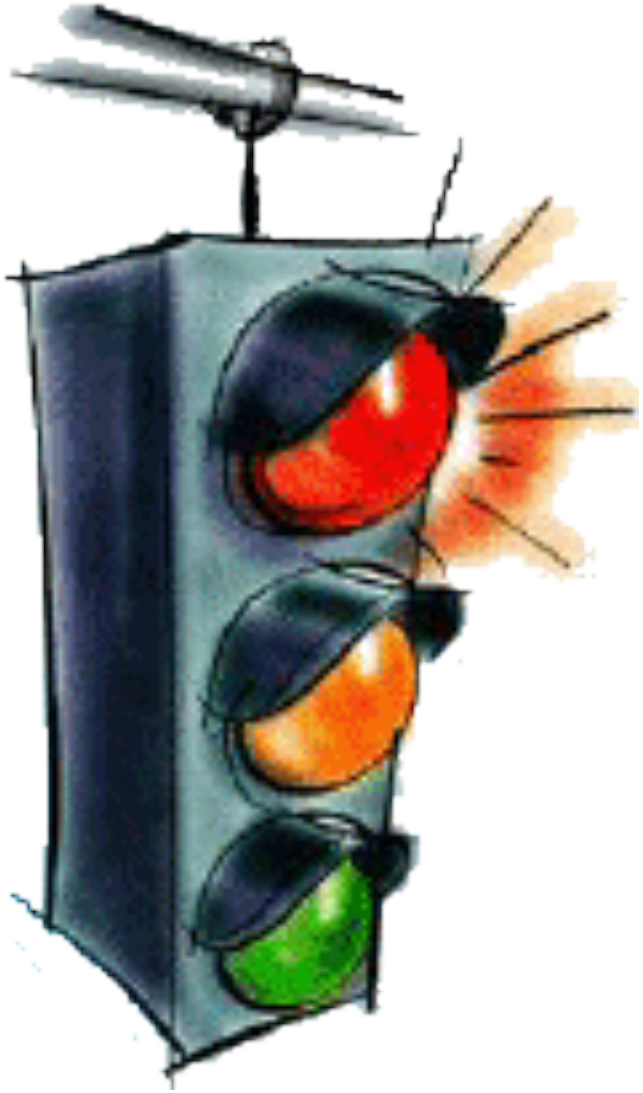
- In the context of colured petri net tokens are typed so that difference instances are always addressed

# Exercise

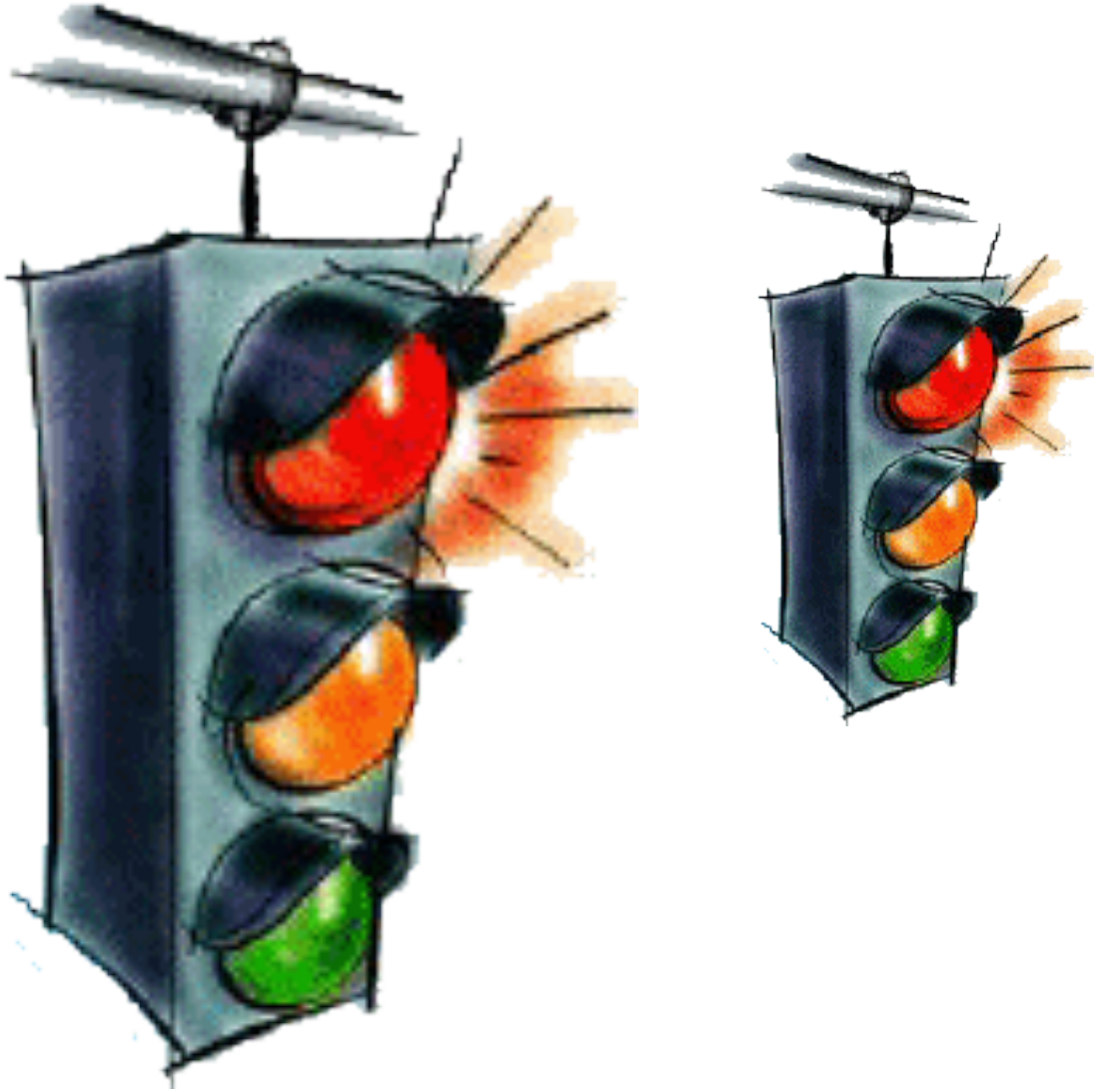
Modelling using petri net the traffic light scenario



# Exercise

Modelling using petri net two traffic light

Complete the net in such a way that the two lights can never be green at the same time

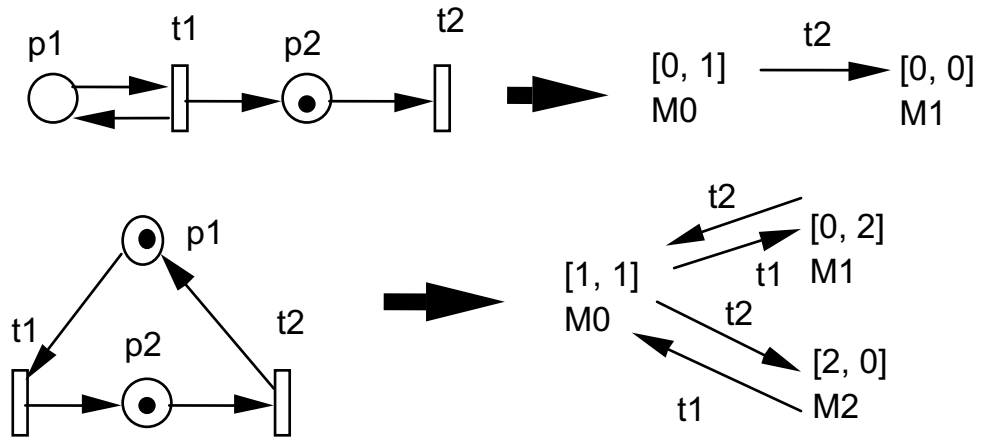


# Reachability tree

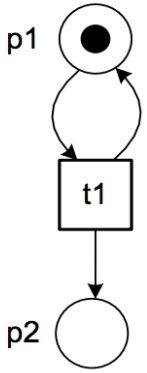
**Definition**: The reachability tree, also called marking graph, of a Petri net (N, M0) is a graph in which

- nodes corresponds to reachable markings
- arcs correpond to feasible transitions

**Remark**: the reachability tree ca be infinite



# Coverability Tree (from Murata)




If, on the path from the root to M, there exists a marking M'' such that M'(p) ≥ M''(p) for each p and M'≠M'' (i.e., M'' is coverable), then replace M'(p) by ω for each p such that M'(p) > M''(p)

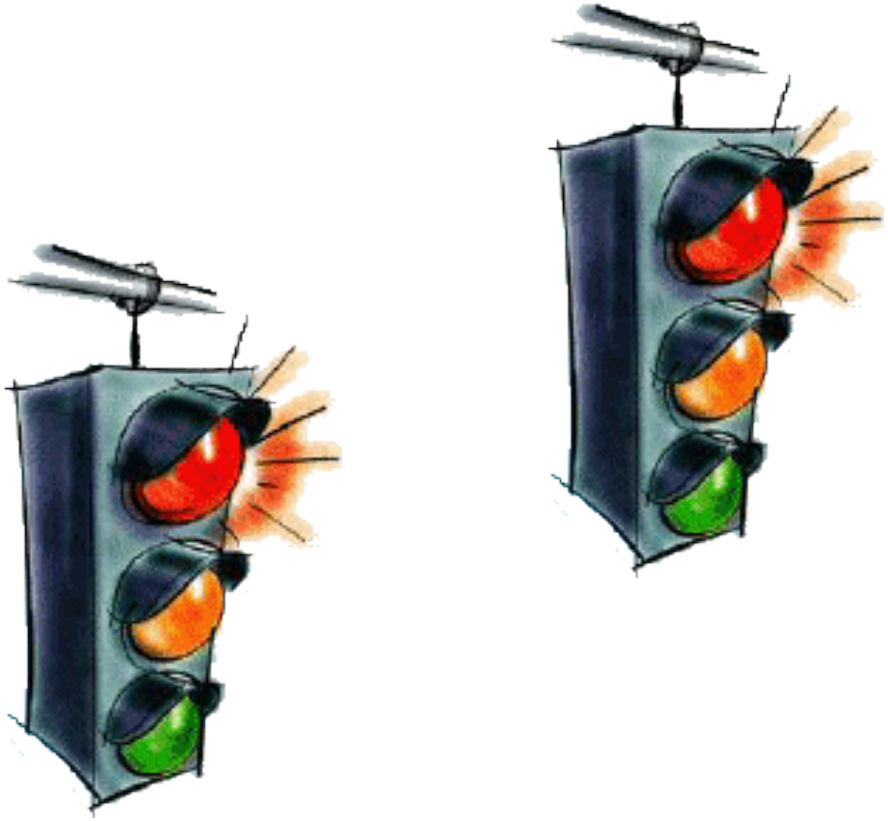
Mo [1,0] ->t1-> M1[1,1]

Mo [1,0] ->t1-> M1[1, ω]

Mo [1,0] ->t1-> M1[1, ω] t1

# Exercise

- Draw the reachability graph of the net




# Mapping from BPMN to Petri Nets

# Mapping a Need for Verification

- Business Processes are typically defined by business experts which ask for graphical and user-friendly notations
- Most notations such as BPMN 2.0 used typically lack precisely defined semantics limiting the possibility of analysis
- In order to formally analysis Business Process we need to derive a formal model of the Business Process being verified
  - From BPMN 2.0 to Petri Net
  - From BPMN 2.0 to Process Algebra
  - From ... To …

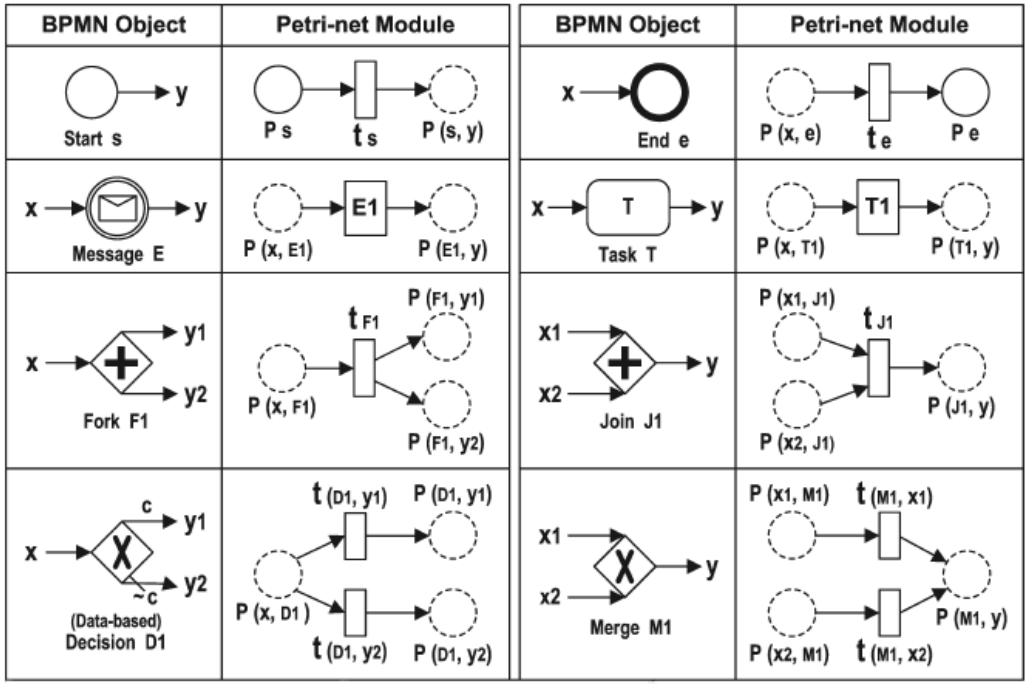
# From BPMN 2.0 to Petri Net (I)

- Raedts et. al, 2007 - *Raedts, I., Petkovic, M., Usenko, Y. S., van der Werf, J. M., Groote, J. F., and Somers, L. J. (2007). Transformation of BPMN models for Behaviour Analysi, In MSVVEIS, pages 126-137.*
  - It converts BPMN 2.0 to Petri Nets extended with inhibitor and reset arcs
  - The proposed mapping is not comprehensive enough in order to clearly address all the BPMN 2.0 elements (i.e. messages, sub-process and loop activity are not included)

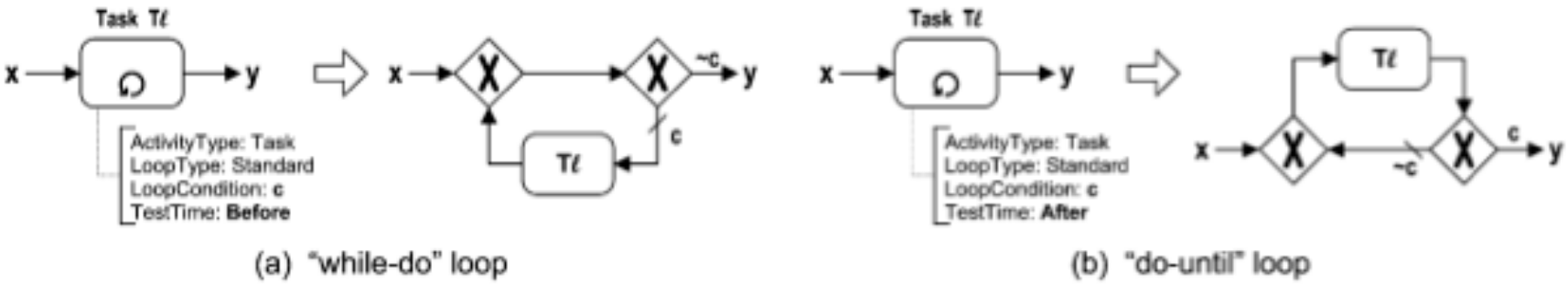
- Awad, 2010 - *Awad, A., Decker, G., and Lohmann, N. (2010). Diagnosing and repairing data anomalies in process models. BPM Workshops, vol. 43 of Lecture Notes in Business Information Processing, pp. 5-16. Springer.*
  - It formalizes the basic data object together with the more traditional control flow elements
  - It is missed a more general and complex mapping on the BPMN 2.0 elements such as sub-process, messages, etc.

# From BPMN 2.0 to Petri Net (II)

- Dijkman, R. M., Dumas, M., and Ouyang, C. (2008). Semantics and analysis of business process models in BPMN. Inf. Softw. Technol., 50(12):1281-1294.
  - Up to our review it the most complete methodology to transform BPMN 2.0 models into Petri
  - It starts from the basic elements and introduces some more details regarding multiple instances, sub-process and exception handling
  - It can be easy extended to the whole specification



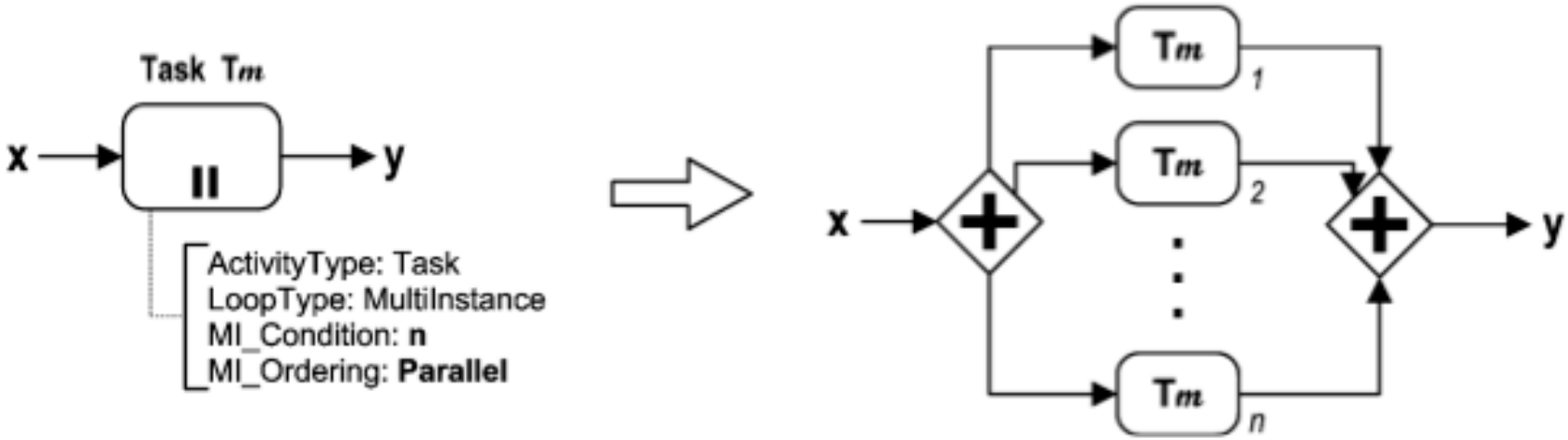
# Macro expansions for repeated activities



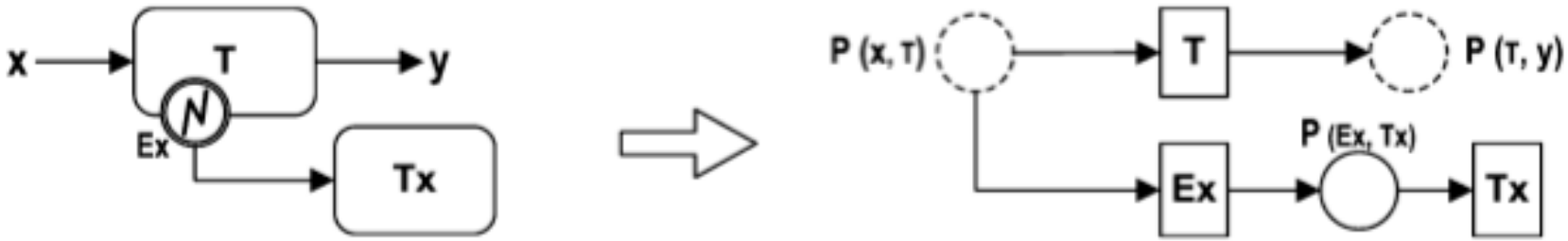

(a) "while-do" loop

(b) "do-until" loop

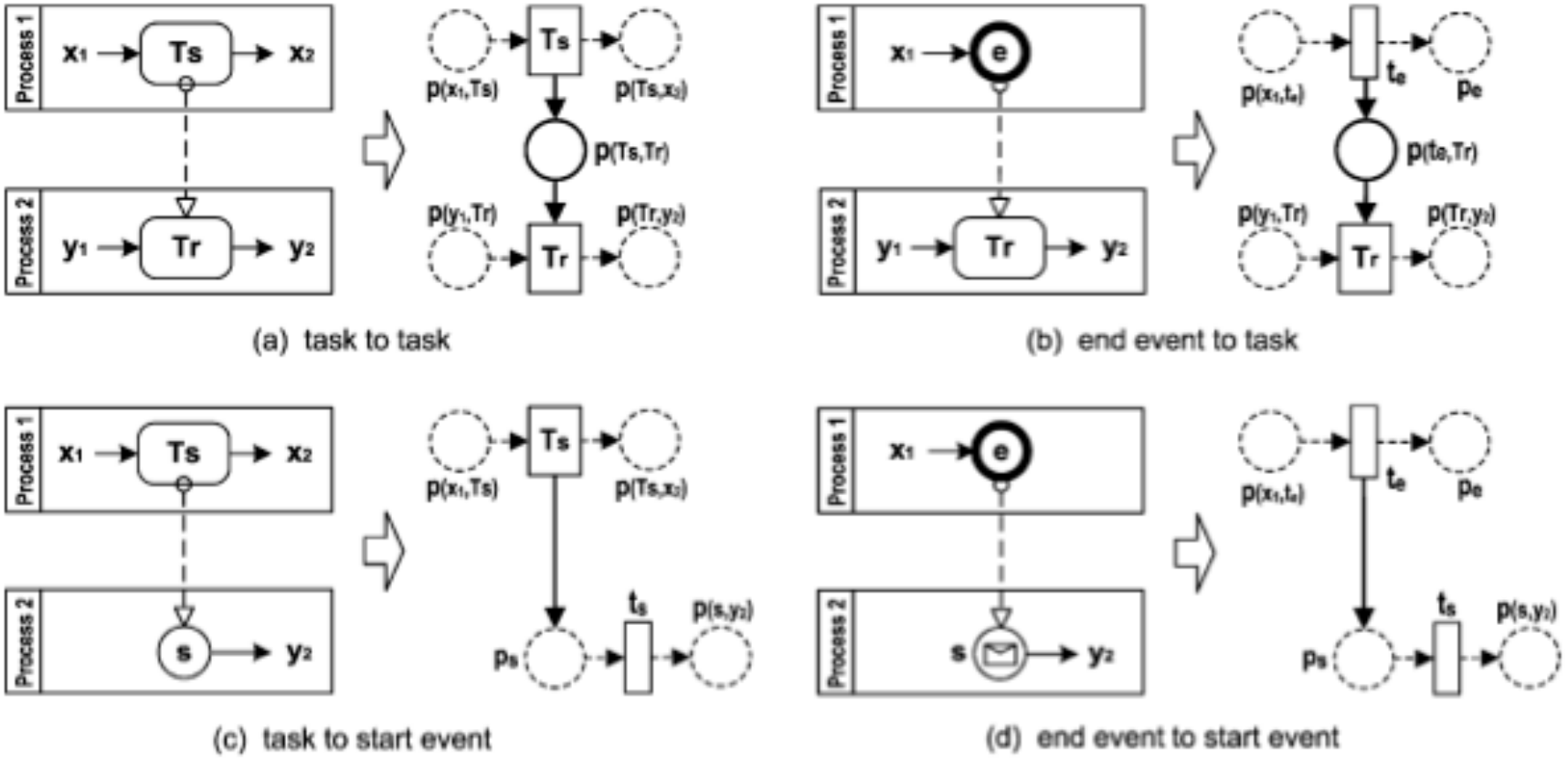
# Macro expansion for a multi-instance activity where n is known at design time



# Mapping of a task with an exception flow



# Mapping of message flows between BPMN processes



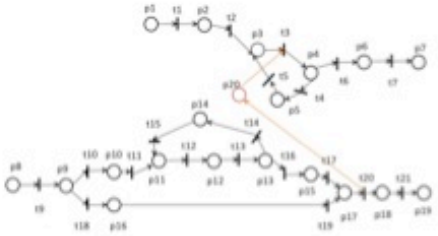

(a) task to task

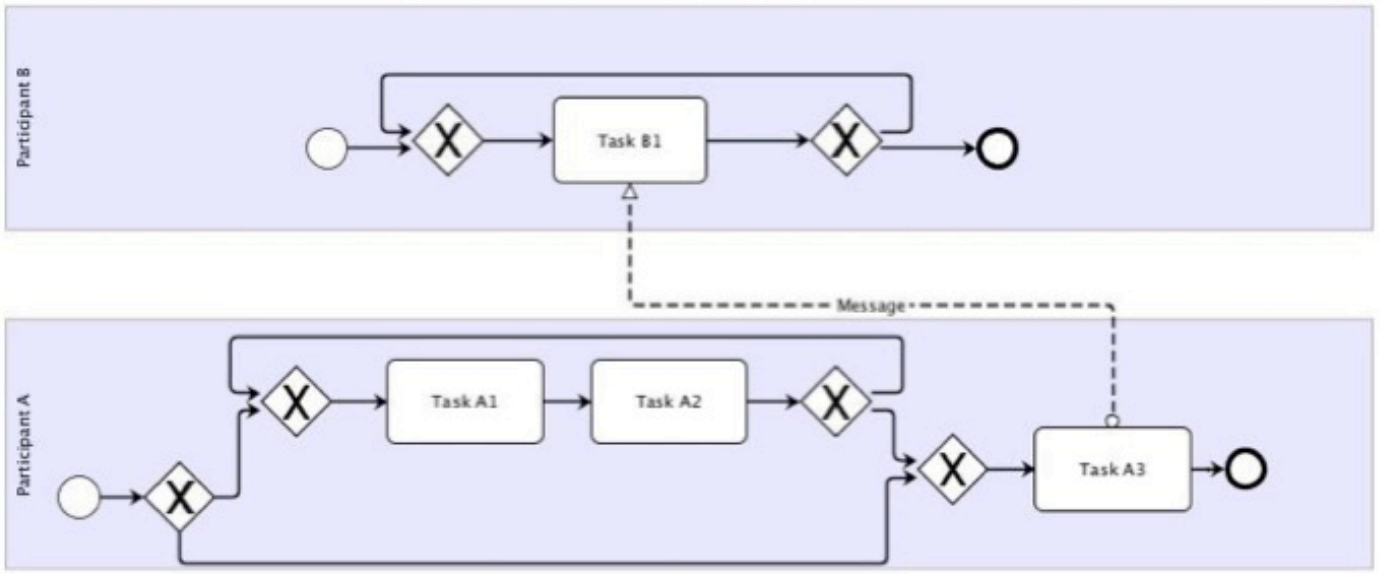
(b) end event to task

(c) task to start event

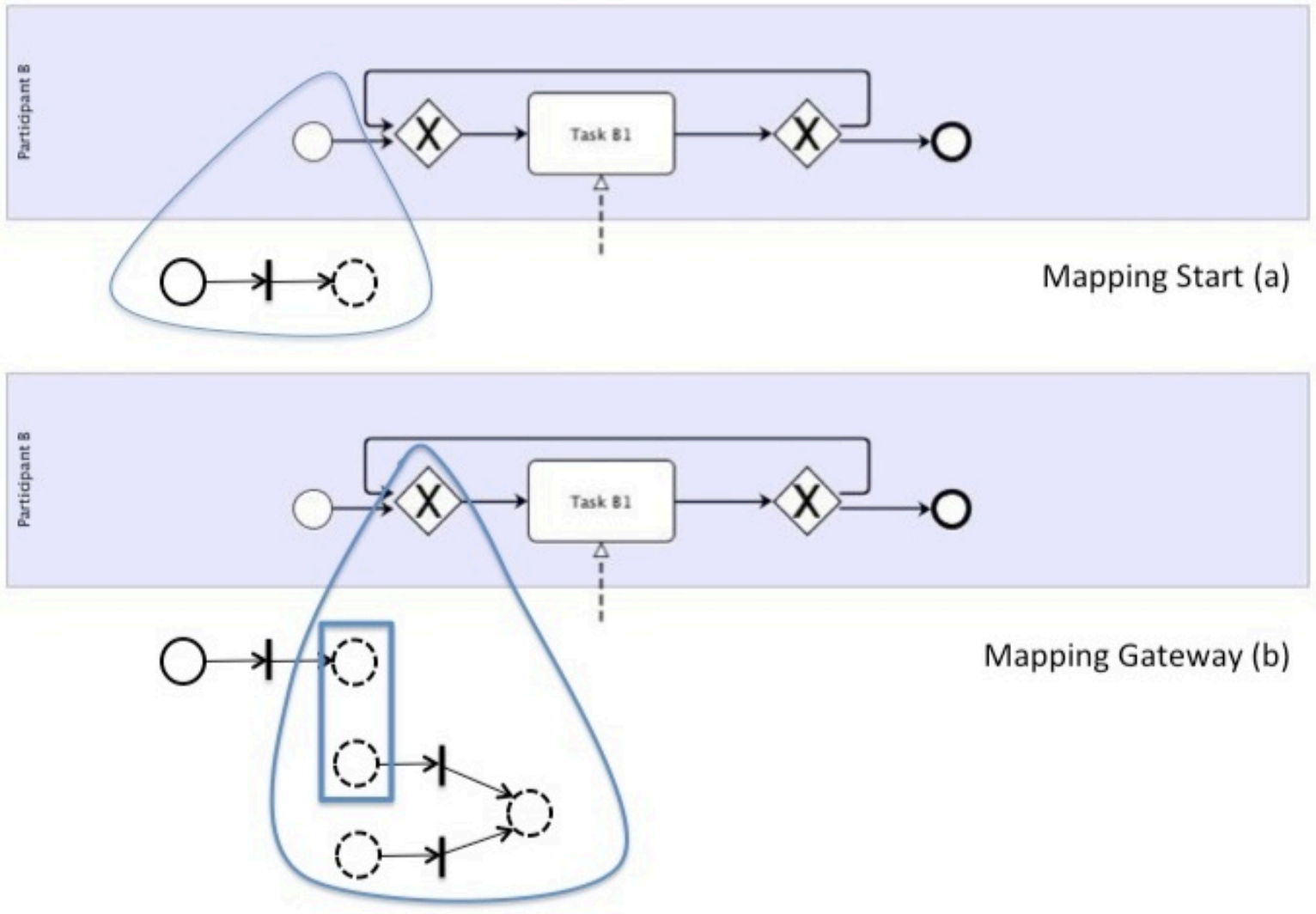
(d) end event to start event

# Basic Example of Mapping BP into Petri Net

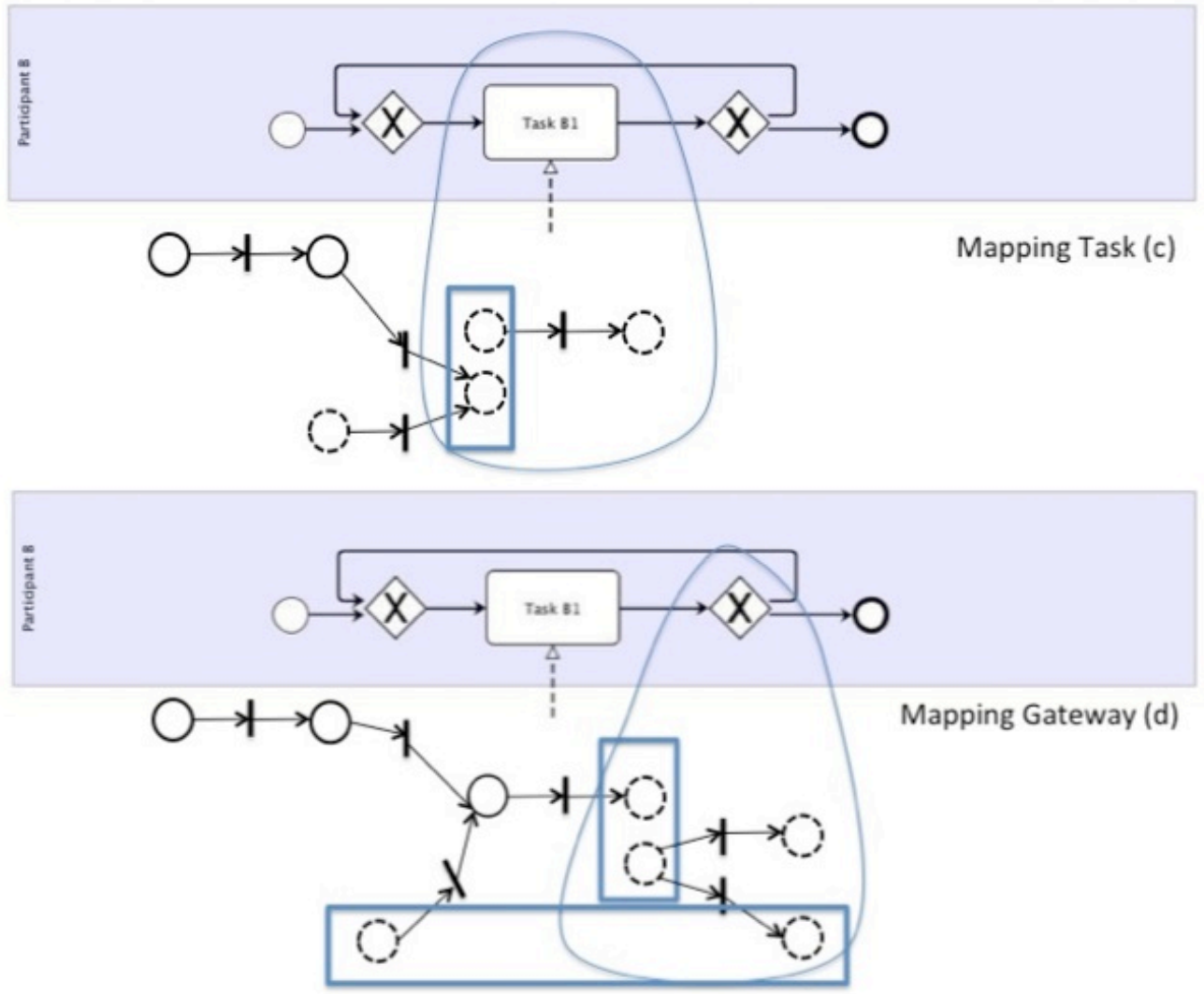




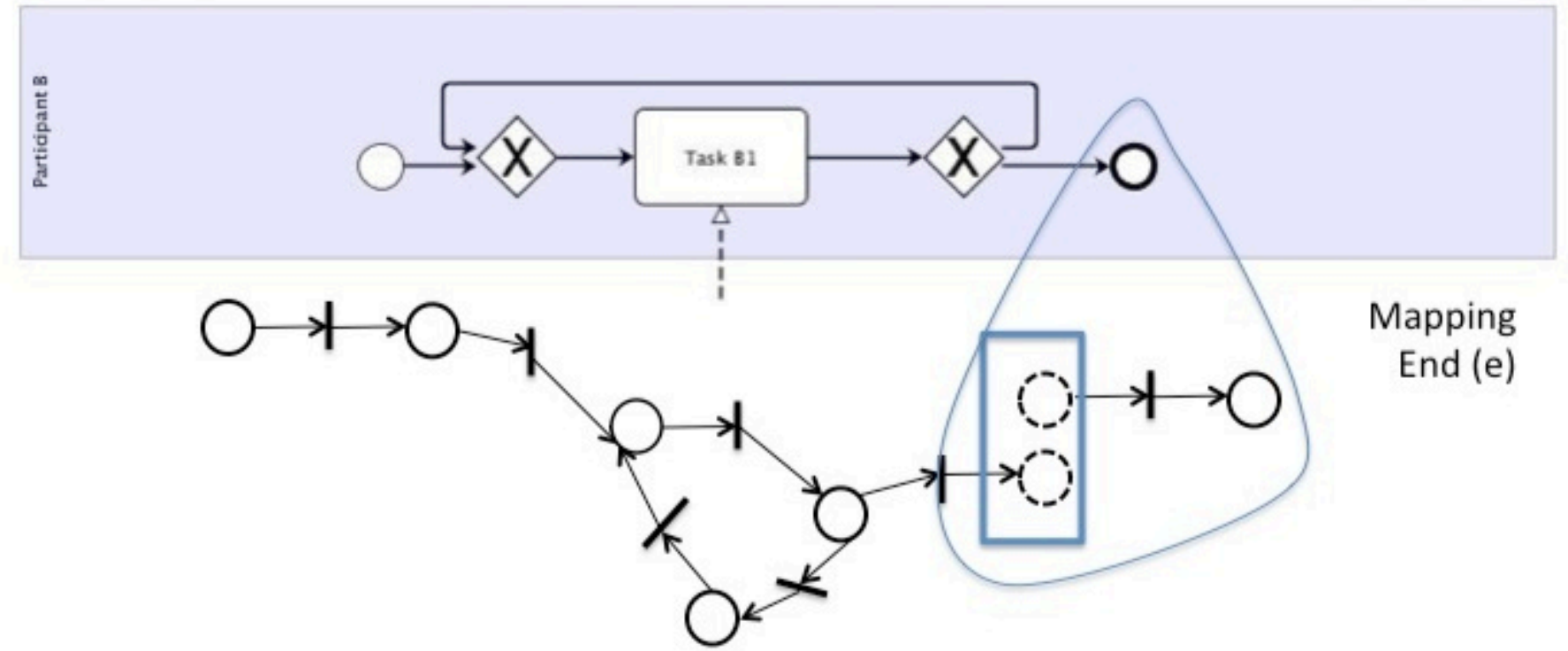
# Basic Example of Mapping BP into Petri Net (step 1 participant B)



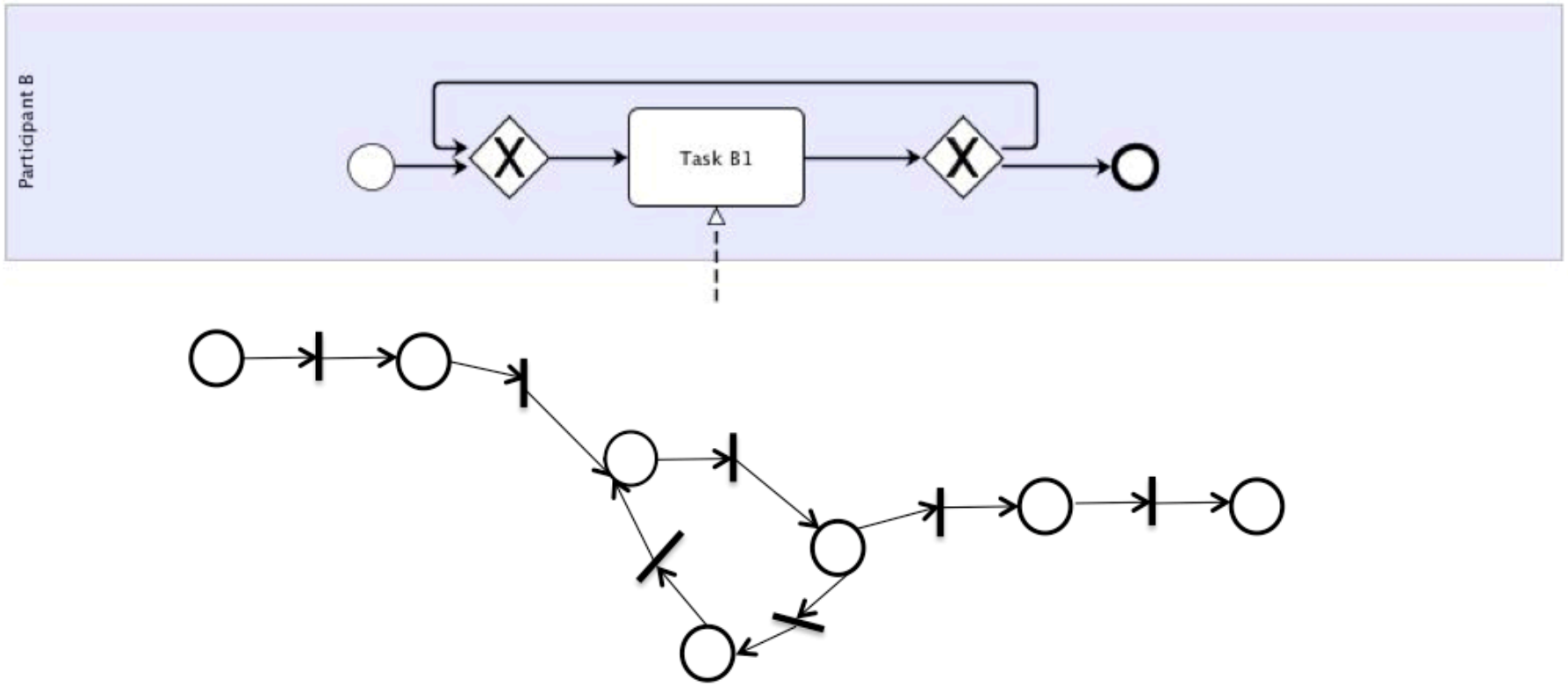
# Basic Example of Mapping BP into Petri Net (step 2 participant B)



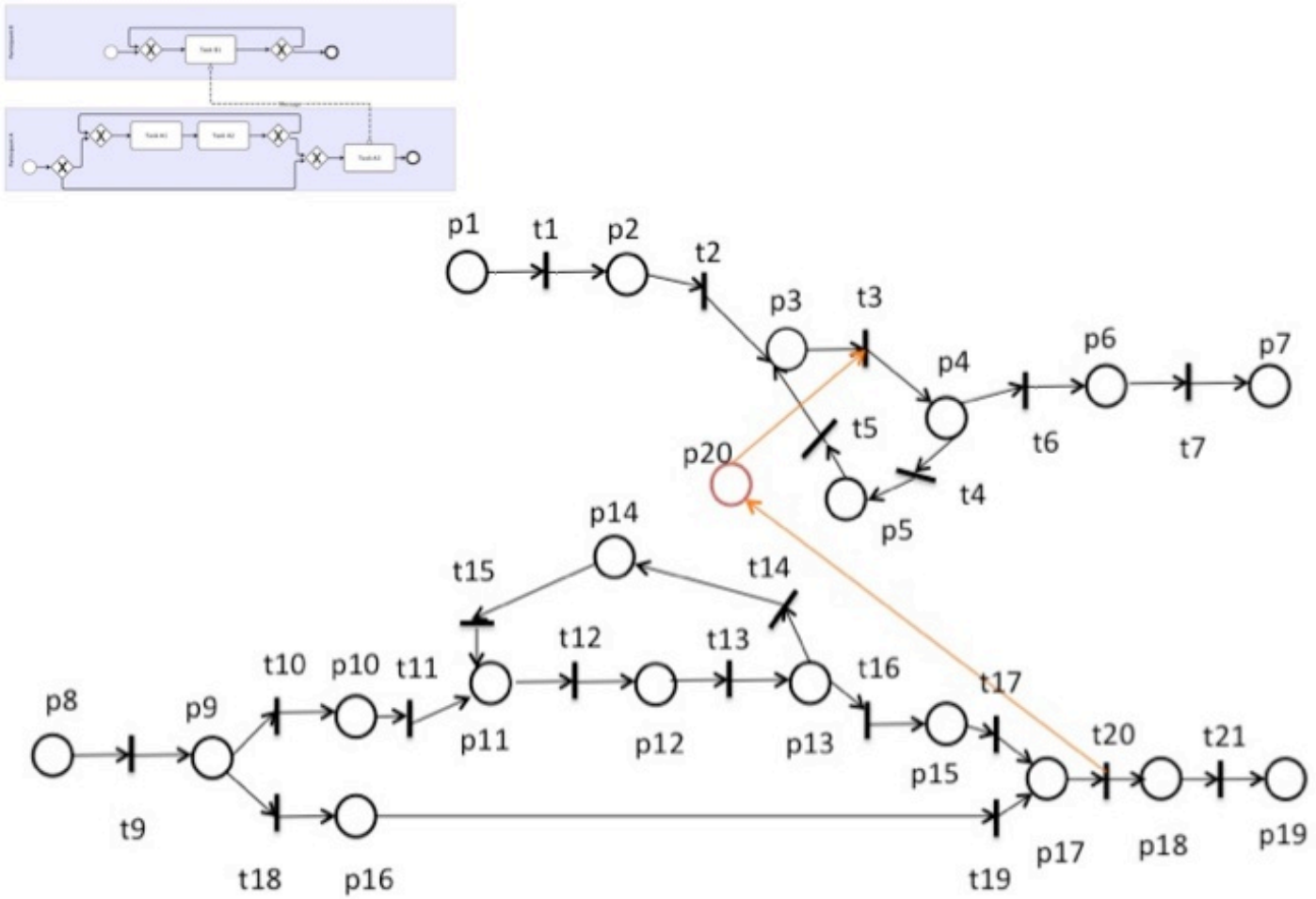
# Basic Example of Mapping BP into Petri Net (step 3 participant B)



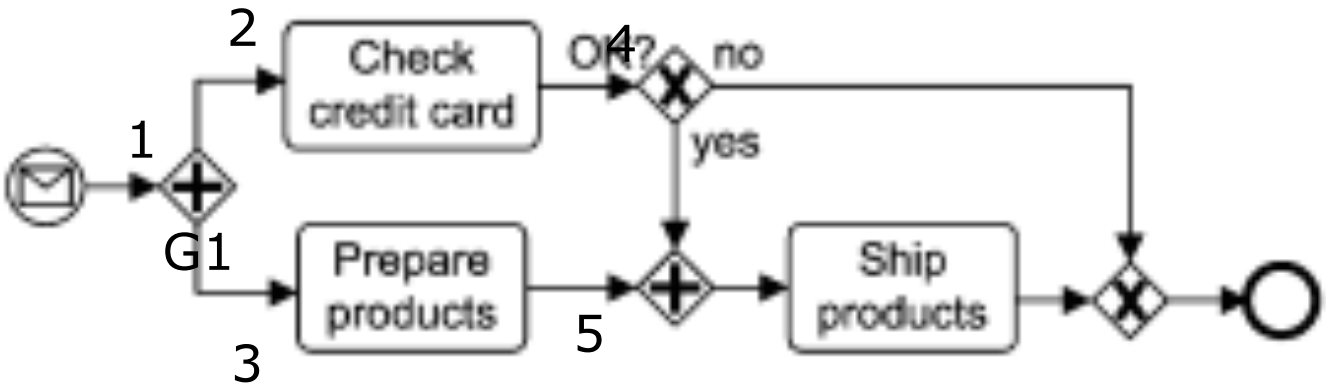
# Basic Example of Mapping BP into Petri Net (participant B)

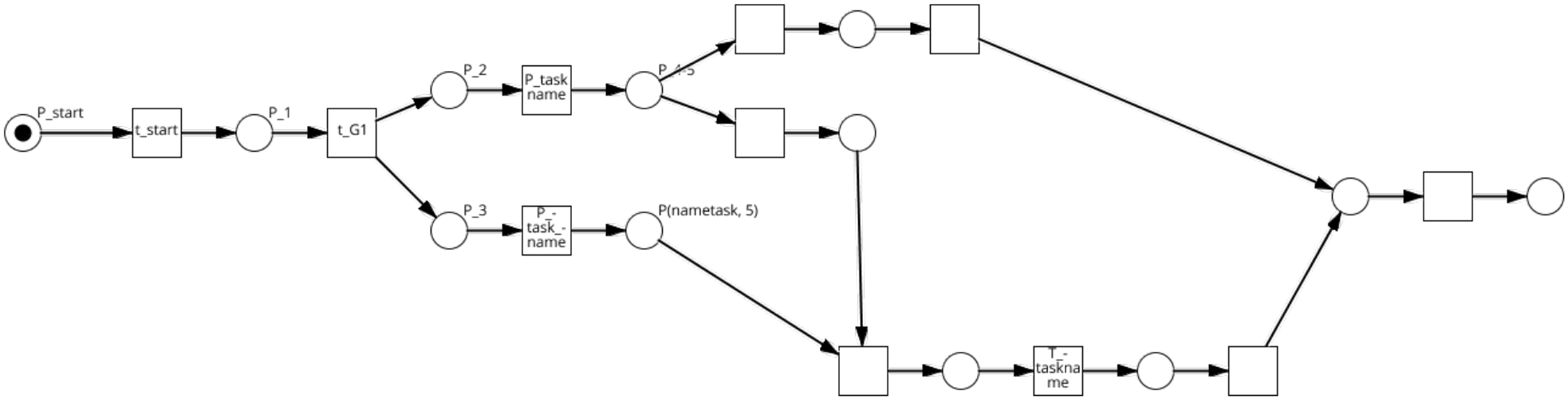


# Basic Example of Mapping BP into Petri Net



# Mapping into Practice



# Mapping into Practice