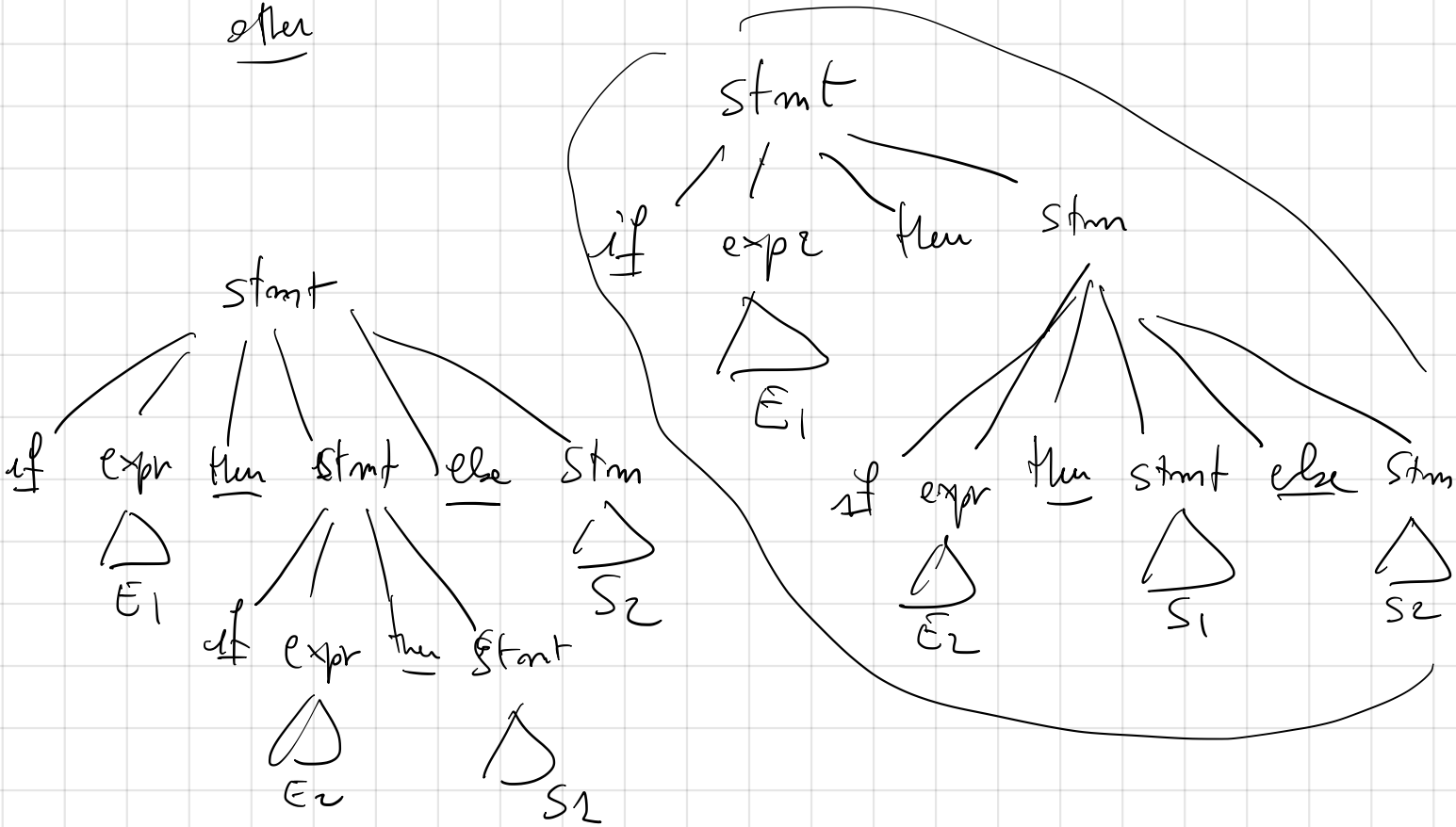


if E_1 then if E_2 then S_1 else S_2

stmt \rightarrow if expr then stmt |

if expr then stmt else stmt |

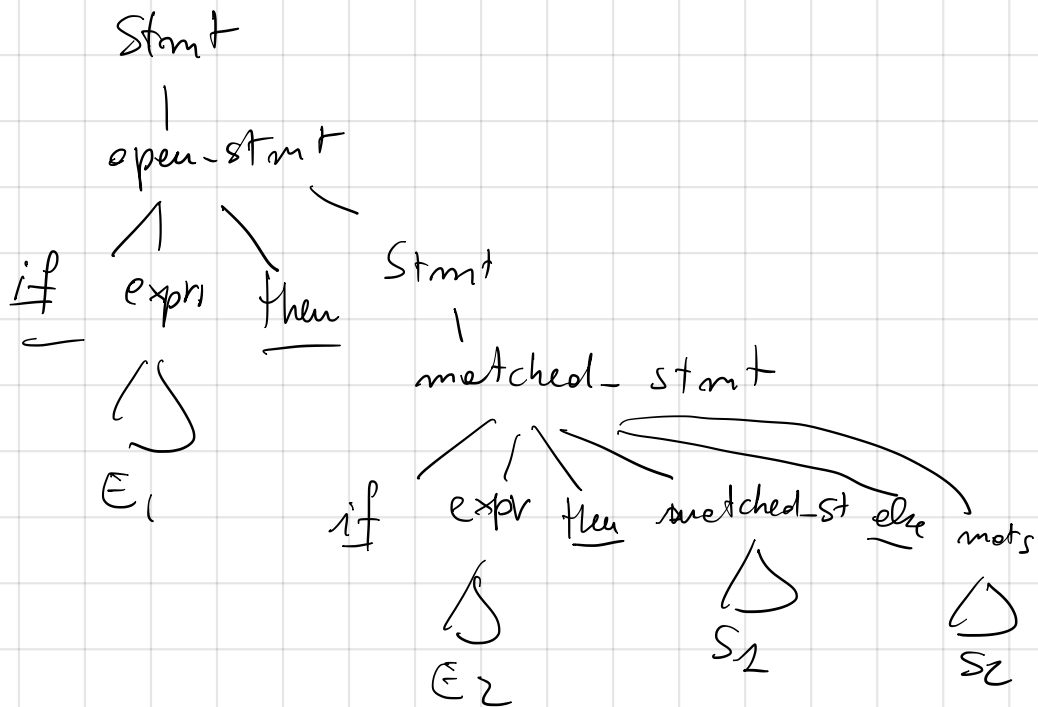
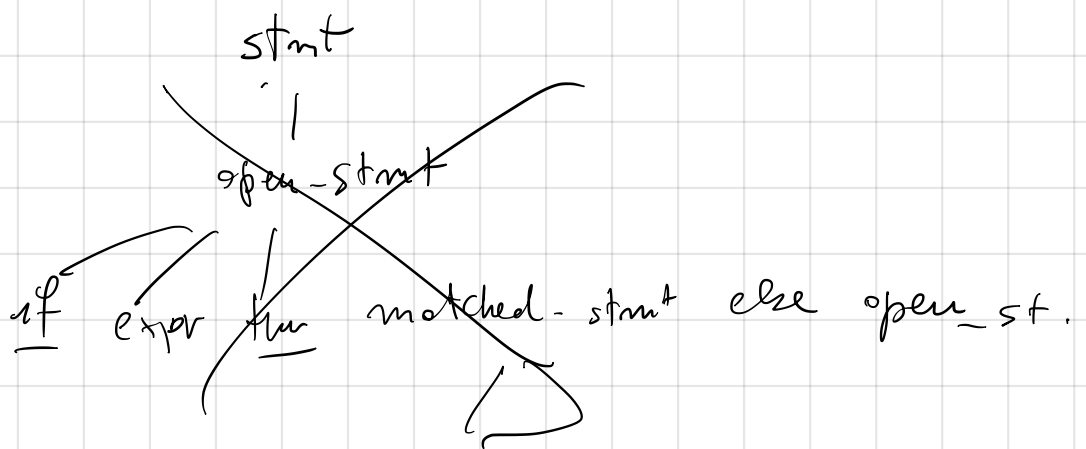
other



stmt \rightarrow matched_stmt | open_stmt

matched_stmt \rightarrow other | if expr then matched_stmt
else matched_stmt

open_stmt \rightarrow if expr then stmt |
if expr then matched_stmt else open_stmt



Regular Expressions

$$\Sigma = \{a, b\}$$

* 1	+	left	S
2	.	left	T
3	*	/	R
4	atoms	/	B

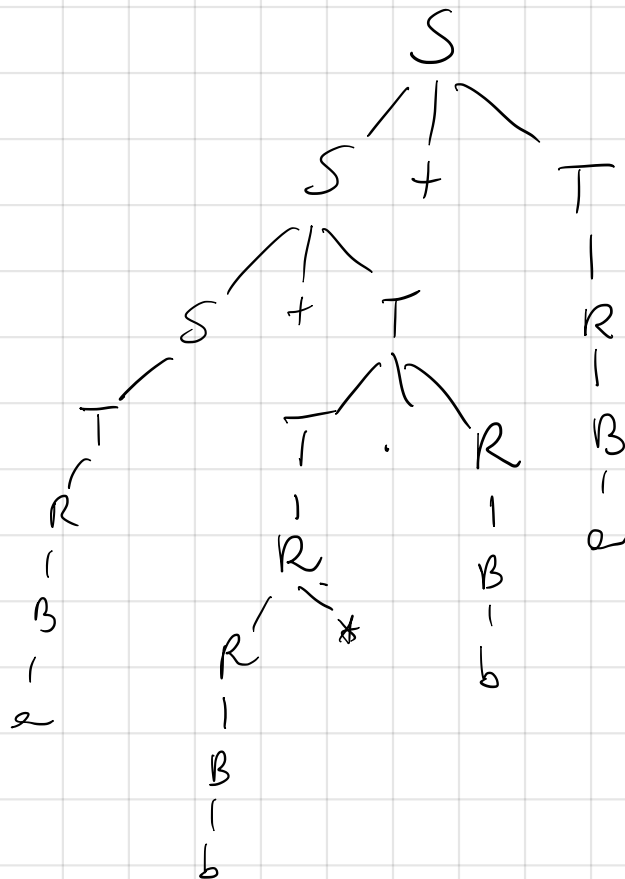
$$S \rightarrow S + T \mid T$$

$$T \rightarrow T \cdot R \mid R$$

$$R \rightarrow R^* \mid B$$

$$B \rightarrow a \mid b \mid \text{eps} \mid (S)$$

$$a + b^* b + a$$

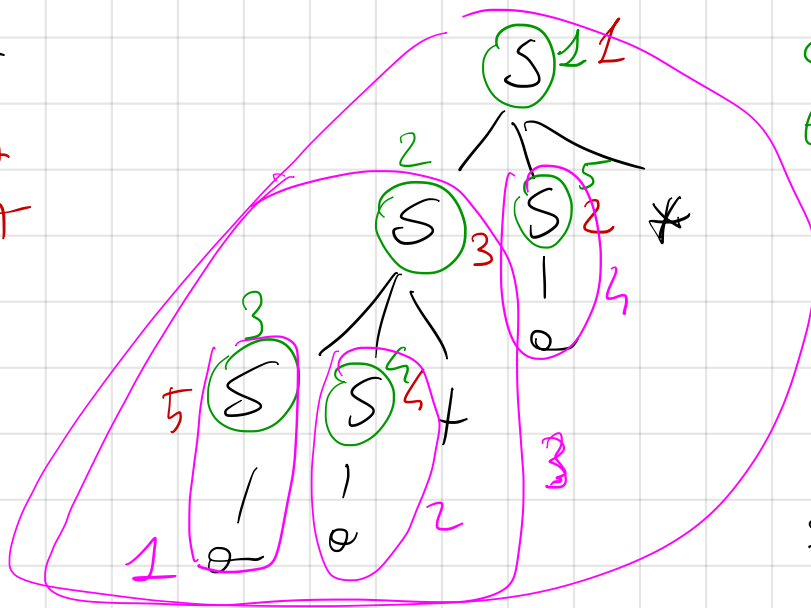


$$S \rightarrow SS+ \mid SS* \mid e$$

ea+e*

depth-first
right-first

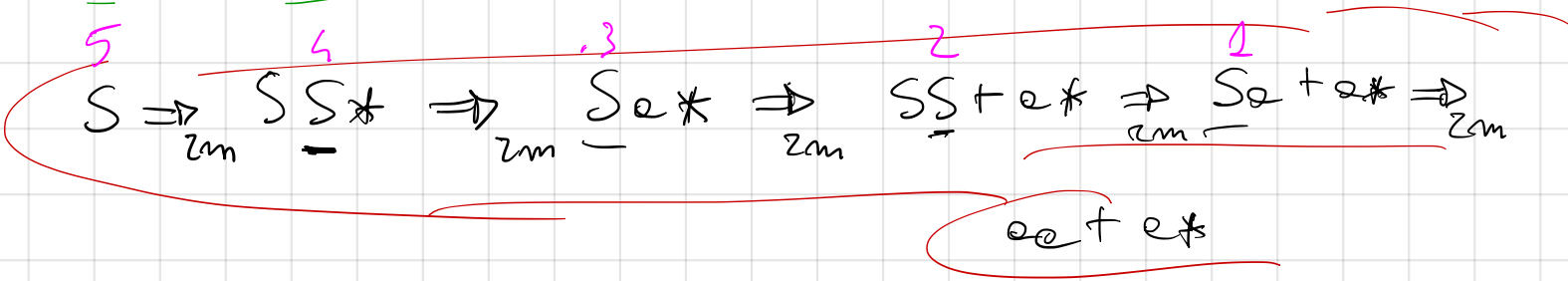
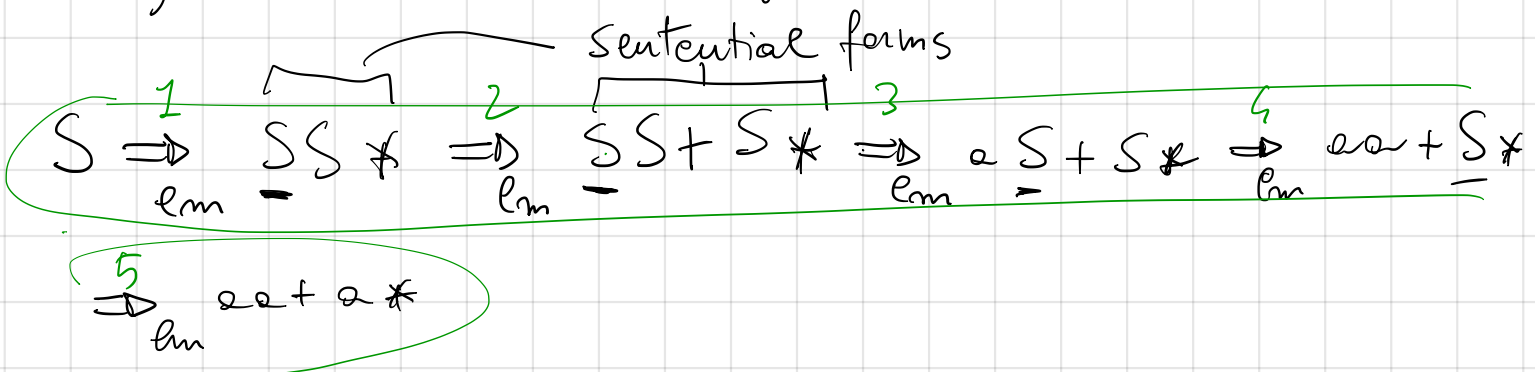
depth-first
left-first



$$7+9*5$$

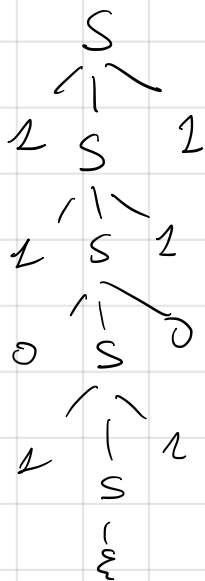
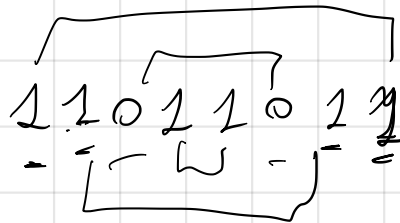
$$79+5*$$

The grammar is not ambiguous



$\mathcal{L} = \{w \in \{0,1\}^* \mid w \text{ is palindromic}\}$

$S \rightarrow \emptyset S \emptyset \mid 1 S 1 \mid \epsilon$

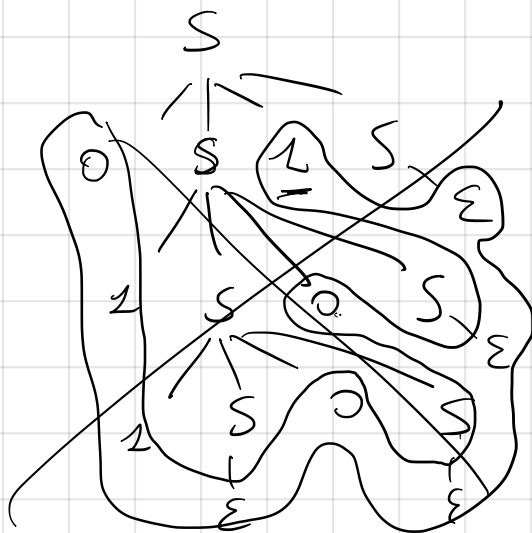


$\mathcal{L} = \{w \in \{0,1\}^* \mid \text{the number of 1's in } w \text{ equals the number of 0's in } w\}$

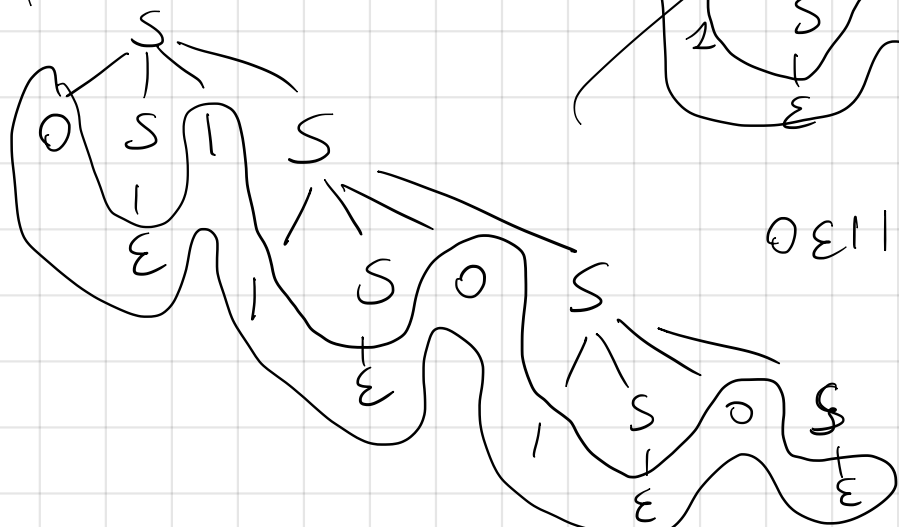
$S \rightarrow \underline{1} S \underline{0} S \mid \underline{0} S \underline{1} S \mid \epsilon$

011010

~~011010 = 3120303110~~



= 3031031130



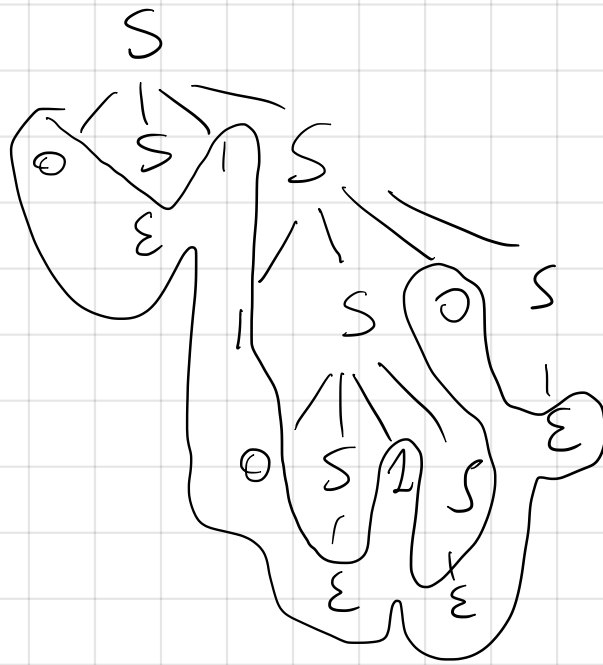
011010
() () ()

011010

()(())

0ε110ε1ε0ε

011010



$L = \{ x \in \{ (,) \}^* \mid x \text{ is a string of balanced parentheses} \}$

$S \rightarrow (S)S \mid \epsilon$

$$L = \{ a^m b^m c^k \mid m=k \text{ OR } m=k, m, k \geq 0 \}$$

$$S \rightarrow S_1 C \mid A S_2$$

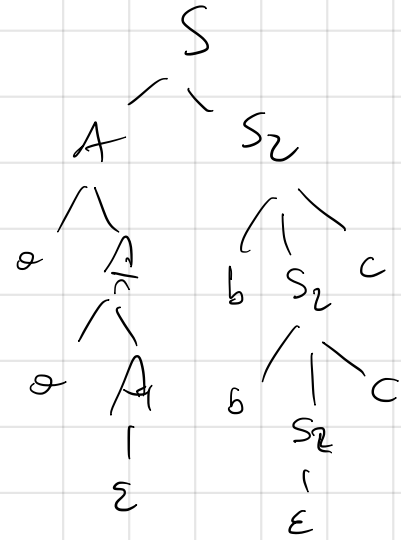
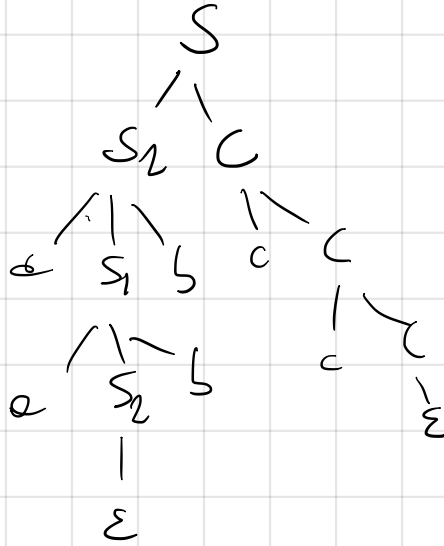
$a b b c c$

$$C \rightarrow c C \mid \epsilon$$

$$A \rightarrow a A \mid \epsilon$$

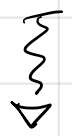
$$S_1 \rightarrow a S_1 b \mid \epsilon$$

$$S_2 \rightarrow b S_2 c \mid \epsilon$$



$$S \rightarrow A_0 | b$$

$$A \rightarrow \underline{A}c | \underline{S}d | \underline{\varepsilon}$$

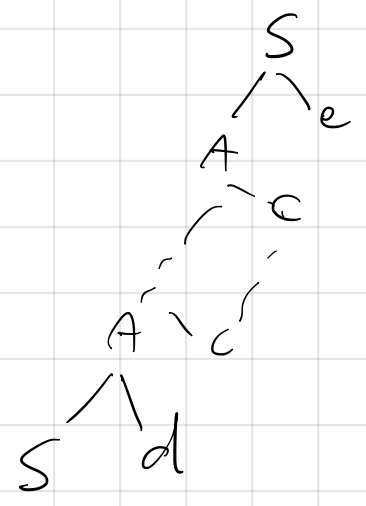
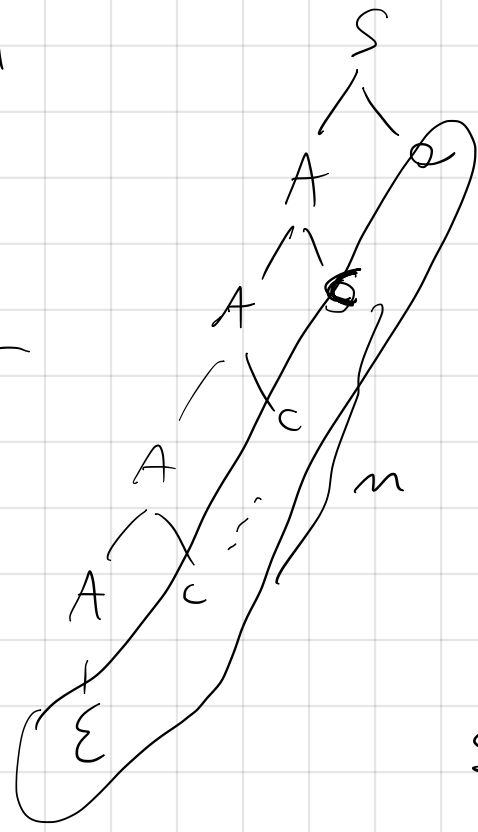
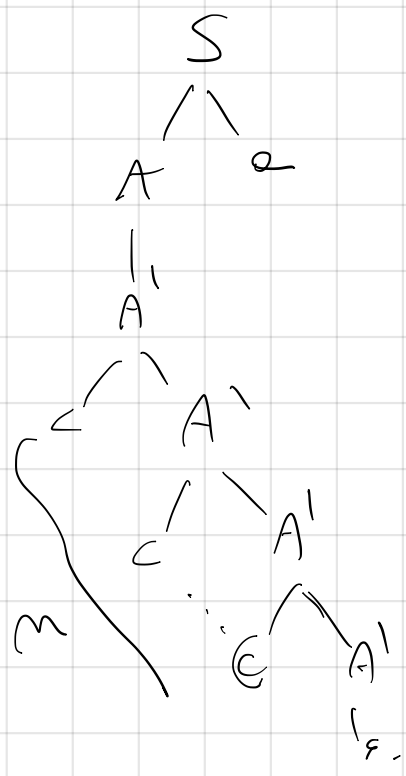


$$\left\{ \begin{array}{l} S \rightarrow A_0 | b \\ A \rightarrow SdA' | \cancel{A'} \\ A' \rightarrow cA' | \varepsilon \end{array} \right.$$

$$A \rightarrow A\alpha | \beta$$



$$\begin{array}{l} A \rightarrow \beta A' \\ A' \rightarrow \alpha A' | \varepsilon \end{array}$$



$$\underline{\underline{S \Rightarrow A_0 \Rightarrow S}}$$

$$\begin{aligned} E &\rightarrow E + T \mid T \\ T &\rightarrow T * F \mid F \\ F &\rightarrow \underline{id} \mid (E) \end{aligned}$$

$$\begin{aligned} E &\rightarrow T E' \\ E' &\rightarrow + T E' \mid \epsilon \\ T &\rightarrow F T' \\ T' &\rightarrow * F T' \mid \epsilon \\ F &\rightarrow \underline{id} \mid (E) \end{aligned}$$

$$\left(\begin{aligned} A &\rightarrow A \alpha \mid B \\ &\Downarrow \\ A &\rightarrow \beta A' \\ A' &\rightarrow \alpha A' \mid \epsilon \end{aligned} \right)$$

$id + id * id$

Abstract
Syntax tree

