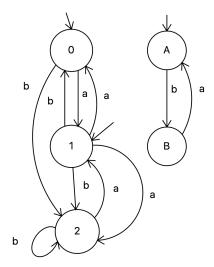
## Master of Science in Computer Science - University of Camerino Reactive Systems Verification A. Y. 2016/2017 Written Test of 12th July 2017 (Appello II)

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## EXERCISE 1 (6 points)

Consider the two following transition systems.



1. Draw the transition system resulting from their product using handshaking with the handshake action set  $H = \{a, b\}$ .

## EXERCISE 2 (9 points)

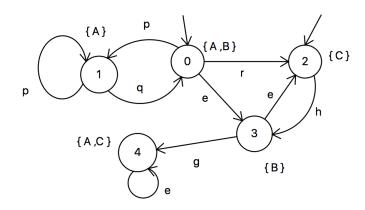
Consider the alphabet  $AP = \{A, B, C, D\}$  and the following linear time properties:

- (a) Whenever A holds then B does not hold for two steps
- (b) B and C hold together only finitely many times
- (c) If C holds infinitely many times then D holds only finitely many times
- (d) C holds at least once and whenever D holds also A must hold

For each property:

- 1. formalise it using set expressions and first order logic;
- 2. formalise it in LTL (you can use the operators next, until, box and diamond, together with all boolean connectives);
- 3. tell if it is a safety, liveness or mixed property; in case it is a safety property provide an NFA for the language of the **minimal** bad prefixes.

Consider the following transition system TS on  $AP = \{A, B, C\}$ .



1. (3 points) Decide, for each LTL formula  $\varphi_i$  below, whether or not  $TS \models \varphi_i$ . Justify your answers! If  $TS \not\models \varphi_i$  provide a path  $\pi \in Paths(TS)$  such that  $\pi \not\models \varphi_i$ .

$$\varphi_1 = \bigcirc \bigcirc A$$

$$\varphi_2 = \Box (A \Rightarrow \Diamond C)$$

$$\varphi_3 = \Diamond \Box (A \lor C)$$

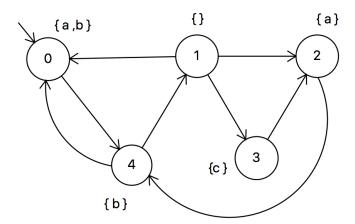
2. Consider the following fairness assumptions:

$$\begin{array}{ll} \psi_1^{\text{fair}} = \{ \{ \}, \{q\}, \{r\} \} & \psi_2^{\text{fair}} = \{ \{ \}, \{r\}, \{q\} \} \\ \psi_3^{\text{fair}} = \{ \{ \}, \{g\}, \{e\} \} & \psi_4^{\text{fair}} = \{ \{ \}, \{ \}, \{g, e\} \} \end{array}$$

- (a) (3 points) Decide whether or not  $TS \models_{\mathrm{fair}} \varphi_2$  under the four different fairness conditions  $\psi^i_{\mathrm{fair}}$ ,  $i \in \{1, 2, 3, 4\}$ , separately. Whenever  $TS \not\models_{\mathrm{fair}} \varphi_2$  provide a path  $\pi \in Paths(TS)$  such that  $\pi \not\models \varphi_1$  and arguing that  $\pi$  is fair with respect to  $\psi^i_{\mathrm{fair}}$ .
- (b) (3 points) Decide whether or not  $TS \models_{\mathrm{fair}} \varphi_3$  under the four different fairness conditions  $\psi^i_{\mathrm{fair}}$ ,  $i \in \{1, 2, 3, 4\}$ , separately. Whenever  $TS \not\models_{\mathrm{fair}} \varphi_3$  provide a path  $\pi \in Paths(TS)$  such that  $\pi \not\models \varphi_6$  and arguing that  $\pi$  is fair with respect to  $\psi^i_{\mathrm{fair}}$ .

## EXERCISE 4 (8 points)

Consider the following transition system



1. Calculate  $\operatorname{Sat}(a \wedge (b \leftrightarrow c))$ ,  $\operatorname{Sat}(\exists (c \vee b)\mathcal{U}a)$  and  $\operatorname{Sat}(\exists \Box (b \vee c))$ . Justify your answers by showing the steps of the algorithm used for the CTL formulas.