

$$\varphi_1 = \Diamond b$$

$T \models \varphi_1$ NO Counterexample: $(S_0 S_1)^\omega = T_1$

$$\psi_{\text{fair}}^1 = \square \Diamond c \rightarrow \square \Diamond b$$

?

$T \models_{\psi_{\text{fair}}^1} \varphi_1$ T_1 violates ψ_{fair}^1 , so it is excluded by fairness

$$\left(\{a, c\} \setminus \{a\} \right)^\omega$$

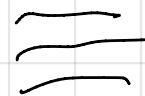
Any other path is fair w.r.t. ψ_{fair}^1

$$so \quad T \models_{\psi_{\text{fair}}^1} \varphi_1 !$$

$$\varphi_5 = (a \vee b) \cup (a \vee c)$$

$$T \models \varphi_5 \text{ on}$$

→ justification:



$T \not\models \varphi_6$ Counterexample is $(S_3 S_4)^\omega \rightsquigarrow (\{b\} \setminus \{a\})^\omega$

$$or \\ S_0 S_2 (S_3 S_4)^\omega \rightsquigarrow \{a, c\} \setminus \{a, b, c\} (\{b\} \setminus \{a\})^\omega$$

$$\varphi_2 \equiv \neg \square \underbrace{(b \rightarrow (b M (e \wedge \neg b)))}$$

PNF $\varphi ::= \text{true} \mid \text{false} \mid \omega \mid \neg a \mid \varphi_1 \wedge \varphi_2 \mid \varphi_1 \vee \varphi_2 \mid \Diamond \varphi \mid$
 $\varphi_1 M \varphi_2 \mid \varphi_1 W \varphi_2$

$$\Box \varphi = \neg \square \neg \varphi$$

$$\begin{array}{c} p \rightarrow q \\ \equiv \end{array}$$

$$\Diamond \varphi = \text{true} M \varphi$$

$$\neg p \vee q$$

$$\varphi_2 \equiv \neg \neg \square \neg (b \rightarrow (b M (e \wedge \neg b)))$$

$$\equiv \Diamond \neg (\neg b \vee (b M (e \wedge \neg b)))$$

$$\equiv \text{true} M \neg (\neg b \vee (b M (e \wedge \neg b)))$$

$$\equiv \text{true} M (\neg \neg b \wedge \neg (b M (e \wedge \neg b)))$$

$$\equiv \text{true} M (b \wedge \neg (e \wedge \neg b) W (\neg b \wedge \neg (e \wedge \neg b)))$$

$$\equiv \text{true} M (b \wedge (\neg e \vee \neg \neg b) W (\neg b \wedge (\neg e \vee \neg \neg b)))$$

$$\equiv \text{true} M (b \wedge (\neg e \vee b) W ((\neg b \wedge \neg e) \vee (\neg b \vee b)))$$

$$\equiv \text{true} M (b \wedge (\neg e \vee b) W ((\neg b \wedge \neg e) \vee \text{true}))$$

$$\equiv \text{true} M (b \wedge (\neg e \vee b) W (\neg b \wedge \neg e))$$