

Programmazione Avanzata

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Corso di Laurea in Informatica (L31) Scuola di Scienze e Tecnologie



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Programming Languages





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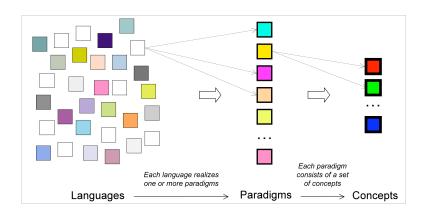
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A programming paradigm is an approach to programming a computer based on a mathematical theory or a coherent set of principles.

Each paradigm supports a set of concepts that makes it the best for a certain kind of problem.

Solving a programming problem requires choosing the right concepts!







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- imperative/procedural: statements are used to change program's state. Imperative programming focuses on describing how a program operates.
- functional: computation is treated as the evaluation of mathematical functions and avoids changing-state and mutable data.
- declarative/logical: expresses the logic of a computation without describing its control flow. A program consists in a set of sentences in logical form, expressing facts and rules about some problem domain.
- object-oriented: it is based on the concept of objects, which may contain both data, the fields, and code, the methods.



In this lecture...



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... we will first introduce basic notions of functional programming...



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...then some basic notions of declarative programming is provided...



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... after that we focus on object-oriented programming...



In this lecture...

... we will first introduce basic notions of functional programming...

...then some basic notions of declarative programming is provided...

... after that we focus on object-oriented programming...

... finally an overview of modern programming languages is provided.



Functional programming in F#: Basic Concepts

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Computations consist in the appropriate compositions of defined functions.

We will consider F#, a modern functional language integrated in the .Net framework.



```
F# (pronounced F sharp)...
... is a strongly typed programming language;
```



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Here we will main consider the functional aspects!

F# programming language Primitive Types (1/2)



- bool, Boolean values (true or false).
- byte, Unsigned byte (from 0 to $2^8 1$).
- sbyte, Signed byte (from -2^7 to $2^7 1$).
- int16, 16-bit integer (from -2^{15} to $2^{15} 1$).
- uint16, 16-bit integer (from 0 to $2^{16} 1$).
- int, 32-bit integer (from -2^{31} to $2^{31} 1$).
- uint32, 32-bit unsigned (from 0 to $2^{32} 1$).
- int64, 64-bit integer (from -2^{63} to $2^{63} 1$).
- \blacksquare uint64, 64-bit unsigned int (from 0 to $2^{64}-1$).
- char, Unicode character values.
- string, Unicode text.
- decimal, Floating point data type that has at least 28 significant digits.

F# programming language Primitive Types (2/2)



- unit, Indicates the absence of an actual value.
- void, Indicates no type or value.
- float32, A 32-bit floating point type.
- float, A 64-bit floating point type.

F# programming language Values (1/2)



- bool: true, false.
- byte, an integer with postfix y (86y).
- sbyte, an integer with postfix uy (86uy).
- int16, an integer with postfix s (86s).
- uint16, an integer with postfix us (86us).
- int, an integer with the optional postfix I (86 or 861).
- uint32, an integer with postfix u or ul (86u or ul).
- int64, an integer with postfix L (86L).
- uint64, an integer with postfix UL (86UL).
- char, a single symbol surrounded by single quotes ('a').

F# programming language Values (2/2)



- string, can be:
 - ... a sequence of characters surrounded by double quotes ("Hello\n\n World!");
 - ... a sequence of characters surrounded by double quotes and prefixed with @ (@"Hello\n\n World!");
 - ... a portion of text (possibly on multiple lines) surrounded by """

```
""" Hello
```

World!"""

- decimal, a floating point value postfixed with M (0.35M).
- unit, the value ().
- float32, a floating point postfixed with f or F (0.35f or 035F).
- float, a floating point in decimal or exponential form (0.35 or 3.5E-1).

F# programming language Basic concepts



Basic construct in F# is let that can be used to associate a name with a value

```
\begin{array}{lll} \text{let num} &=& 10 \\ \text{let str} &=& \text{"F\#"} \end{array}
```



Basic concepts

Basic construct in F# is let that can be used to associate a name with a value

```
let num = 10
let str = "F#"
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Each name has a type that is inferred from the associated expression!

F# programming language Basic concepts



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```
let num = 10
let str = "F#"
```

Each name has a type that is inferred from the associated expression!

Above:

- num has type int;
- str has type string;

F# programming language Operators



Arithmetic Operators: +, -, *, /, %, **;

Comparison Operators: =, <, <=, >, >=, <>;

Boolean Operators: not, ||, &&;

Bitwise Operators: &&&, |||, ^^^, ~~~, <<<, >>>;

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Comparison Operators:
$$=$$
, $<$, $<=$, $>$, $>=$, $<>$;

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Arithmetic and Comparison operators are overloaded: the exact type depends on the type of their argument!

Differently from Java, no implicit cast is done!

F# programming language Simple type errors!



```
let x = 86u //x has type ubyte
let y = 86 //y has type int
let z = x+y //This is an error!!!!
```

F# programming language Basic concepts



Functions are first-class values and can be associated with names as any other built-in types:



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let
$$f1 = fun \times -> x+1$$

let $f2(x) = x+1$

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$$f3(x,f) = f(x+2)+1$$

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Function types have the form: type1 -> type2



The idea of type inference is that you do not have to specify the types of F# constructs except when the compiler cannot conclusively deduce the type.



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Type annotations can be used to help the compiler to infer the expected type.



```
Type inference
```

```
//No annotation, inferred type: int*int \rightarrow int let f(x,y) = x+y
```



Type inference

```
//No annotation , inferred type: int*int \rightarrow int let f(x,y)=x+y
//Parameter x is annotated as float , inferred type: float* float \rightarrow float let f(x): float , y) = x+y
```



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Type inference
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//No annotation, inferred type: int*int \rightarrow int let f(x,y) = x+y

//Parameter x is annotated as float, inferred type: float* float \rightarrow float let f(x): float, f(x): float, f(x): float, f(x): float, f(x): float, f(x): float f(
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Type inference

```
//No annotation, inferred type: int*int -> int
let f(x,y) = x+y
//Parameter x is annotated as float, inferred type: float*
 float -> float
let f(x: float, y) = x+y
//Name x is annotated as float, inferred type: float*float
 -> float
let f(x,y) = (x: float)+y
//Return type of f is float,
// inferred type: float *float -> float
let f(x,y): float = x+y
```



Partial evaluation

Let us consider the following functions:

$$let f1(x,y) = x+y$$

$$\mbox{let} \ \ \mbox{f2} \ \ \mbox{x} \ \ \mbox{y} \ = \ \mbox{x+y}$$



Partial evaluation

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val f2 :
$$x:int \rightarrow y:int \rightarrow int$$



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The two approaches are in fact equivalent! The second one is the standard (and more efficient).



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let select x y z w = if (x=y) then z else w
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The following type is inferred for function select:

```
val select : x:'a -> y:'a -> z:'b -> w:'b -> 'b
    when 'a : equality
```



Recursive functions...

In functional programming the use of recursive definition is crucial.



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```
 \begin{array}{lll} \text{let fib x} &=& \\ \text{if x<1 then} \\ &\text{1} \\ &\text{else} \\ &\text{(fib x-1)+(fib x-2)} \\ \end{array}
```



Recursive functions...

In functional programming the use of recursive definition is crucial.

```
 \begin{array}{lll} \text{let fib } & x & = \\ & \text{if } x < 1 \text{ then} \\ & 1 \\ & \text{else} \\ & \text{(fib } x - 1) + \text{(fib } x - 2) \\ \end{array}
```

This definition is not correct! The symbol fib is not defined when the body of the function is evaluated!



Recursive functions

In functional programming the use of recursive definition is crucial.

```
let fib x =
  if x<1 then
   1
  else
     (fib x-1)+(fib x-2)</pre>
```

This definition is not correct! The symbol fib is not defined when the body of the function is evaluated!

```
let rec fib(x) = //Note here the use of 'rec' if x<=2 then 1 else (fib x-1)+(fib x-2)
```

F# programming language Tuples...



A tuple is a grouping of unnamed but ordered values, possibly of different types.

```
(element, ..., element)
```



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Example:

Tuples...

```
let fib(x) =
  let rec _fib(x) =
    if x<=2 then
      (1,1)
  else
      let (a,b)=_fib(x-1)
      in
            (a+b,a)
  in
    let (a, _) = _fib(x)
    in
      a</pre>
```



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You can define a list by explicitly listing out the elements, separated by semicolons and enclosed in square brackets;

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let list123 = [ 1; 2; 3 ] //Type int list let emptylist = [] //Type 'a list!
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let list1 = [1..10]
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let list1 = [1..10]
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List operations:

- :: is used to add an element at the beginning of the list: a:: list1
- @ is used to concatenate two lists: |1@|2



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// Match expression.
match test-expression with
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Conditions are boolean expressions that can be used to limit the selection.



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$$a_n x^n + \cdots + a_1 x + a_0$$



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let poly =
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Write a function eval that received in input a list of coefficients and a value \times computes the value of the polynomial.



Solution 1:

```
let rec eval clist (x: float) =
   match clist with
   | [] -> 0.0
   | c::tail -> c*(x**float(clist.Length-1))+(eval tail x)
```



Solution 1:

```
let rec eval clist (x: float) =
   match clist with
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Solution 2:



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Find the first element in a list matching a predicate:



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Find the first element in a list matching a predicate:

The type of findFirstMatching is:

```
pred:('a -> bool) -> l:'a list -> 'a option
```

Excercises



Ex. 0: Download and install F# developing environment. See instructions available here:

https://docs.microsoft.com/en-us/dotnet/fsharp/get-started/

Ex. 1: Write a function that given in input a and b computes their mcd.

Ex. 2: Write a function that given in input *n* returns true if *n* is a *prime* number and false otherwise.

Ex. 3: Write a function that given in input an integer *n* computes the list of its prime factors.