

# Functional programming in F#: Data Structures

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# Summary of previous lectures

In the previous lecture we have...:

- introduced basic principles of **programming paradigms**;
- introduced basic concept of a **functional language**, F#:
  - elementary types;
  - expressions;
  - function definitions;
  - type inference;
  - lists.



**Ex. 1:** Write a function that given in input  $a$  and  $b$  computes their *mcd*.

## Exercises: solutions

**Ex. 1:** Write a function that given in input  $a$  and  $b$  computes their *mcd*.

```
let mcd (a: uint32) (b: uint32) =  
  let rec _mcd x y =  
    if y=0u then x  
    else _mcd y (x%y)  
  in  
    _mcd (max a b) (min a b)
```

## Exercises: solutions

**Ex. 2:** Write a function that given in input  $n$  returns true if  $n$  is a *prime number* and false otherwise:

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```
let isPrime n =  
  let limit = int(sqrt(float(n)))+1  
  let rec test n x =  
    if x>=limit then true  
    else if n%x = 0 then false  
    else test n (x+1)  
  in  
    test n 2
```

## Exercises: solutions

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```

let primeFactors n =
  let rec divP x y =
    if x%y <> 0 then (x,0)
    else
      let (z,p) = divP (x/y) y
      in
        (z,p+1)
  let rec combP x y =
    if x<2 then []
    else
      match divP x y with
      | (_,0) -> combP x (y+1)
      | (z,p) -> y::(combP z (y+1))
  in
    combP n 2
  
```



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```
let mutable counter = 0  
  
let step () = x <- x+1;x
```

## Loop expressions. . .

The `for ... to` expression is used to iterate in a loop over a range of values of a loop variable:

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for identifier = start [ to | downto ] finish do  
    body-expression
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```

The `while ... do` expression is used to perform iterative execution (looping) while a specified test condition is true:

```
while test-expression do  
    body-expression
```



## Examples: list of prime factors. . .

```

let primeFactors2 n =
  let compP x y =
    let mutable count = 0
    let mutable v = x
    while v%y = 0 do
      v <- v/y
      count <- count+1
    (v, count)
  let mutable v = n
  let mutable count = 2
  let mutable res = []
  while v >= 2 do
    let (r,p) = divP v count
    v <- r;
    if p>0 then res <- count::res;
    count <- count+1;
  res

```

## Custom data type: Records. . .

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```
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    [ mutable ] label2 : type2;  
    ...  
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### Example:

```
type MyPoint = { x: float ; y: float }
```

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    [ mutable ] label2 : type2;  
    ...  
}
```

### Example:

```
type MyPoint = { x: float ; y: float }
```

You can initialize records by using the labels that are defined in the record:

```
let p = { x=10.0 ; y=10.0 }
```

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```
let distance p1 p2 =  
    ((p1.x-p2.x)**2.0+(p1.y-p2.y)**2.0)**0.5
```

## Custom data type: Records...

Fields in a record are accessible via the standard `name.field` notation:

```
let distance p1 p2 =
    ((p1.x-p2.x)**2.0+(p1.y-p2.y)**2.0)**0.5
```

Records can be used with pattern matching. You can specify some fields explicitly and provide variables for other fields:

```
let pointInfo p =
  match p with
  | { x=0.0 ; y=0.0 } ->
    printf "This is the origin!"
  | { x=0.0 ; y=_ } ->
    printf "This point is located on the x-axes!"
  | { x=_ ; y=0.0 } ->
    printf "This point is located on the y-axes!"
  | { x=xval ; y=yval } ->
    printf "This point is located at (%f,%f)" xval yval
```



## Custom data type: Discriminated Unions...

Discriminated unions provide support for values that can be one of a number of named cases, possibly each with different values and types:

```
type type-name =  
  | case-identifier1 [of [ fieldname1 : ] type1 [ * [  
    fieldname2 : ] type2 ... ]  
  | case-identifier2 [of [fieldname3 : ] type3 [ * [  
    fieldname4 : ] type4 ... ]
```

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  fieldname2 : ] type2 ...]
  | case-identifier2 [of [fieldname3 : ] type3 [ * [
  fieldname4 : ] type4 ...]
```

### Example:

```
type Shape =
  | Rectangle of width : float * length : float
  | Circle of radius : float
  | Prism of width : float * float * height : float
```

## Example: Binary Search Trees!

Binary search trees keep their keys in sorted order:

- elements are inserted/removed from the tree by following the principle of binary search;
- elements traverse the tree from root to leaf by making decisions on the base of comparison.

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Binary search trees keep their keys in sorted order:

- elements are inserted/removed from the tree by following the principle of binary search;
- elements traverse the tree from root to leaf by making decisions on the base of comparison.

## Exercise:

1. develop a data type for BST;
2. implement basic operations on BST...
  - insertion;
  - search;
  - deletion.

# Example: Binary Search Trees!

## Data type:

```
type bstree =  
  EMPTY  
  | NODE of value: int * left: bstree * right: bstree
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type bstree =  
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```

## Add a value in the tree:

```
let rec add v t =  
  match t with  
  | EMPTY -> NODE(v,EMPTY,EMPTY)  
  | NODE(v1,l,r) when v1<v -> NODE(v1,l,add v r)  
  | NODE(v1,l,r) -> NODE(v1,add v l,r)
```

# Example: Binary Search Trees!

## Search for an element:

```
let rec contains v t =  
  match t with  
  | EMPTY -> false  
  | NODE(v1, -, -) when v1=v -> true  
  | NODE(v1, l, r) when v1<v -> contains v r  
  | NODE(v1, l, r) -> contains v l
```

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## Search for an element:

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```

## Merging trees:

```

let rec merge t1 t2 =
  match t1, t2 with
  | EMPTY, _ -> t2
  | _, EMPTY -> t1
  | NODE(v1, l1, r1), NODE(v2, l2, r2) when v1<v2 ->
      NODE(v1, l1, merge r1 t2)
  | NODE(v1, l1, r1), NODE(v2, l2, r2) ->
      NODE(v2, l2, merge r2 t1)
  
```



# Example: Binary Search Trees!

## Removing an element:

```
let rec remove v t =  
  match t with  
  | EMPTY -> EMPTY  
  | NODE(v1, l, r) when v1=v -> merge l r  
  | NODE(v1, l, r) when v1<v -> NODE(v1, l, remove v r)  
  | NODE(v1, l, r) -> NODE(v1, remove v l, r)
```

## Exercises. . .

**Ex. 4** Implement function `size` that given a tree `t` computes the number of elements stored in `t`.

**Ex. 5** Implement function `height` that given a tree `t` computes its height (an empty BST has height equal to 0).

**Ex. 6** Implement function `balance` that given a tree `t` computes a tree `t1` with the same elements its height (an empty BST has height equal to 0).

**Ex. 7** Implement AVL data structure.

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The exact type of elements in a `bstree` could be chosen by the programmer!

**We can use Generics!**

## Generics. . .

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In F# function values, methods, properties, and aggregate types such as classes, records, and discriminated unions can be generic.

# Generics...



## Generics. . .

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Generic functions and types enable you to write code that works with a variety of types without repeating the code for each type:

```
/ Explicitly generic function.  
let function-name<type-parameters> parameter-list =  
function-body  
  
// Explicitly generic type.  
type type-name<type-parameters> type-definition
```

# Generics. . .

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### Example:

```
let max a b = if a > b then a else b
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This function has type 'a -> 'a -> 'a when 'a comparison.



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### Example:

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let max a b = if a > b then a else b
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This function has type 'a -> 'a -> 'a when 'a comparison.

Above when 'a comparison is a **constraint**.

## Example...

We can change the definition of `bstree` as follows:

```
type bstree <'T when 'T: comparison> =  
    EMPTY  
    | NODE of value: 'T * left: 'T bstree * right: 'T bstree
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```

We have not to change the functions `add`, `contains`, and `remove`!